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Article

A Game-Theoretic Analysis of *Baccara Chemin de Fer*

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Abstract: Assuming that cards are dealt with replacement from a single deck and that each of Player and Banker sees the total of his own two-card hand but not its composition, *baccara* is a 2×2^{88} matrix game, which was solved by Kemeny and Snell in 1957. Assuming that cards are dealt without replacement from a d -deck shoe and that Banker sees the composition of his own two-card hand while Player sees only his own total, *baccara* is a 2×2^{484} matrix game, which was solved by Downton and Lockwood in 1975 for $d = 1, 2, \dots, 8$. Assuming that cards are dealt without replacement from a d -deck shoe and that each of Player and Banker sees the composition of his own two-card hand, *baccara* is a $2^5 \times 2^{484}$ matrix game, which is solved herein for every positive integer d .

Keywords: *baccara*; *chemin de fer*; sampling without replacement; matrix game; strict dominance; kernel; solution

Classification: MSC primary 91A05; secondary 91A60

1. Introduction

The game of *baccara chemin de fer* (briefly, *baccara*) played a key role in the development of game theory. Bertrand's [1] analysis of whether Player should draw or stand on a two-card total of 5 was the starting point of Borel's investigation of strategic games [2]. Borel [3,4] described Bertrand's study

as “extremely incomplete” but did not himself contribute to *baccara*. It is unfortunate that Borel was unaware of Dormoy’s [5] work, which was less incomplete. Von Neumann [6], after proving the minimax theorem, remarked that he would analyze *baccara* in a subsequent paper. But a solution of the game would have to wait until the dawn of the computer age. Kemeny and Snell [7], assuming that cards are dealt with replacement from a single deck and that each of Player and Banker sees the total of his own two-card hand but not its composition, found the unique solution of the resulting 2×2^{88} matrix game. In practice, cards are dealt without replacement from a *sabot*, or shoe, containing six 52-card decks mixed together. Downton and Lockwood [8], allowing a d -deck shoe dealt without replacement and assuming that Banker sees the composition of his own two-card hand while Player sees only his own total, found the unique solution of the resulting 2×2^{484} matrix game for $d = 1, 2, \dots, 8$. They used an algorithm of Foster [9].

Our aim in this paper is to solve the game without simplifying assumptions. We allow a d -deck shoe dealt without replacement and allow each of Player and Banker to see the composition of his own two-card hand, making *baccara* a $2^5 \times 2^{484}$ matrix game. We derive optimal Player and Banker strategies and determine the value of the game, doing so for every positive integer d . We too make use of Foster’s [9] algorithm. We suspect that these optimal strategies are uniquely optimal, but we do not have a proof of uniqueness.

It will be convenient for what follows to classify the game-theoretic models of *baccara* in two ways. First, we classify them according to how the cards are dealt.

- Model A. Cards are dealt with replacement from a single deck.
- Model B. Cards are dealt without replacement from a d -deck shoe.

Second, we classify them according to the information available to Player and Banker about their own two-card hands.

- Model 1. Each of Player and Banker sees the total of his own two-card hand but not its composition.
- Model 2. Banker sees the composition of his own two-card hand while Player sees only his own total.
- Model 3. Each of Player and Banker sees the composition of his own two-card hand.

(We do not consider the fourth possibility.) Thus, Model A1 is the model of Kemeny and Snell [7], Model B2 is the model of Downton and Lockwood [8], and Model B3 is our primary focus here. Model A2 was discussed by Downton and Holder [10], but Models A3, B1, and B3 have not been considered before, as far as we know.

The analysis under Model B applies exactly to the first round, and it also applies exactly to any subsequent round provided neither Player nor Banker retains knowledge of the cards dealt during previous rounds. This is a consequence of the exchangeability of a well-shuffled shoe.

Like others before us, we restrict our attention to the classical parlor game of *baccara chemin de fer*, in contrast to the modern casino game. (Following Deloche and Oguer [11], we use the authentic French spelling “baccara” rather than the more conventional “baccarat” to emphasize this.) The rules are as follows. The role of Banker rotates among the players (counter-clockwise), changing hands after a Banker loss or when Banker chooses to relinquish his role. Banker announces the amount he is willing

to risk, and the total amount bet on Player's hand cannot exceed that amount. After a Banker win, all winnings must be added to the bank unless Banker chooses to withdraw. Denominations A, 2–9, 10, J, Q, K have values 1, 2–9, 0, 0, 0, 0, respectively, and suits are irrelevant. The total of a hand, comprising two or three cards, is the sum of the values of the cards, modulo 10. In other words, only the final digit of the sum is used to evaluate a hand. Two cards are dealt face down to Player and two face down to Banker, and each looks only at his own hand. The object of the game is to have the higher total (closer to 9) at the end of play. Winning bets on Player's hand are paid by Banker at even odds. Losing bets on Player's hand are collected by Banker. Hands of equal total result in a tie or a *push* (no money is exchanged). A two-card total of 8 or 9 is a *natural*. If either hand is a natural, play ends. If neither hand is a natural, Player then has the option of drawing a third card (but this option is heavily constrained; see below). If he exercises this option, his third card is dealt face up. Next, Banker, observing Player's third card, if any, has the option of drawing a third card. This completes the game. Since several players can bet on Player's hand, Player's strategy is restricted. He must draw on a two-card total of 4 or less and stand on a two-card total of 6 or 7. When his two-card total is 5, he is free to draw or stand as he chooses. (The decision is usually made by the player with the largest bet.) It is safe to say that the constraints on Player's strategy are in his best interest.¹ Banker, on whose hand no one can bet, has no constraints on his strategy.

In the modern casino game, not only is Banker's strategy highly constrained but the casino collects a five percent commission on Banker wins. That makes the game a bimatrix game, which is not considered here.

In Section 3, we show how to evaluate the payoff matrix. We emphasize Model B3 but treat the other models as well. In Section 4, we use strict dominance to reduce the payoff matrix under Model B3 to $2^5 \times 2^{n_d}$, where n_d depends on the number of decks d and satisfies $18 \leq n_d \leq 23$. We get similar reductions of the other models. To proceed further, in Section 5 we re-examine the unique solution of Kemeny and Snell [7] under Model A1 and notice that there are multiple solutions under Models A2 and A3. In Section 6 we derive the unique solution under Model B1 for every positive integer d . Model B1 is of interest because it shows the price of the “with replacement” assumption more clearly than do Models B2 and B3. In Section 7 we re-derive the unique solution of Downton and Lockwood [8] under Model B2, extending it to every positive integer d . These results lead us in Section 8 to a solution under Model B3 for every positive integer d . The feature of the game that allows this is that, under Model B3, the kernel is 2×2 .² The two Banker pure strategies specified by the kernel are dependent on d , while the

¹ Technically, this has not been proved. But the usual argument is that, if Player's strategy were unconstrained, then, under Model A1, when drawing on a two-card total of i , the conditional distribution of Player's three-card total, given that his third card is not 0-valued, is uniform over the nine other totals; in particular, the three-card total is lower than the two-card total with conditional probability $i/9$, otherwise it is higher. This strongly suggests that Player should draw on two-card totals of 0–3 and stand on two-card totals of 6 and 7. Suppose we allow Player the freedom to draw or stand on two-card totals of 4 and 5. Then we have a $2^2 \times 2^{88}$ matrix game, which can be reduced to $2^2 \times 2^{11}$ using Lemma 1 below. We can then check that the two Player pure strategies in which he stands on a two-card total of 4 are strictly dominated. See Problem 5.16 of [12].

² Quoting [13], the Shapley–Snow theorem [14] “states that corresponding to a basic optimal strategy for either player in any matrix game there is a square submatrix (the kernel) whose associated game has the same value as the given game

two Player pure strategies specified by the kernel are independent of d . All computations were carried out in infinite precision using *Mathematica*.

2. Preliminaries

In this section we present two lemmas that will be key to our analysis. They concern the strategic form of a certain matrix game, which is best described in extensive form. Player I chooses a pure strategy $i \in \{0, 1, \dots, m - 1\}$, where $m \geq 2$. Then nature chooses the value of Z_i , a random variable with distribution $P(Z_i = l) = p_i(l)$ for $l = 0, 1, \dots, n$, where $n \geq 1$. If $Z_i = 0$, play ends and player I’s expected gain is $a_i(0)$. If $Z_i = l \in \{1, 2, \dots, n\}$, then player II observes l (but not i) and based on this information chooses a “move” $j \in \{0, 1\}$. If player II chooses move 1 (respectively, move 0), player I’s expected gain is $a_{i,1}(l)$ (respectively, $a_{i,0}(l)$).

If Player II could observe not only $l (\neq 0)$ but i as well, then he would simply choose the smaller of $a_{i,1}(l)$ and $a_{i,0}(l)$. Lemma 1 implies that if, having observed l , player II would make the same choice for every i , then he should make that choice knowing only l . Lemma 2 provides a solution method in the special case of Lemma 1 in which $m = 2$ based on the idea that player II’s best response to each mixed strategy of player I can be readily evaluated.

We can convert the game to strategic form by noticing that player II’s pure strategies can be identified with subsets $T \subset \{1, 2, \dots, n\}$, with player II choosing move 1 if $Z_i \in T$ and move 0 if $Z_i \in T^c$. The following lemma was used implicitly in [7–9] and explicitly in [12].

Lemma 1. *Let $m \geq 2$ and $n \geq 1$ and consider an $m \times 2^n$ matrix game of the following form. Player I has m pure strategies, labeled $0, 1, \dots, m - 1$. Player II has 2^n pure strategies, labeled by the subsets $T \subset \{1, 2, \dots, n\}$. For $i = 0, 1, \dots, m - 1$, there exist probabilities $p_i(0) \geq 0, p_i(1) > 0, \dots, p_i(n) > 0$ with $p_i(0) + p_i(1) + \dots + p_i(n) = 1$ together with a real number $a_i(0)$, and for $l = 1, 2, \dots, n$, there exists a real $m \times 2$ matrix*

$$\begin{pmatrix} a_{0,0}(l) & a_{0,1}(l) \\ a_{1,0}(l) & a_{1,1}(l) \\ \vdots & \vdots \\ a_{m-1,0}(l) & a_{m-1,1}(l) \end{pmatrix}$$

The $m \times 2^n$ matrix game has payoff matrix with (i, T) entry given by

$$a_{i,T} := p_i(0)a_i(0) + \sum_{l \in T} p_i(l)a_{i,1}(l) + \sum_{l \in T^c} p_i(l)a_{i,0}(l) \tag{1}$$

for $i \in \{0, 1, \dots, m - 1\}$ and $T \subset \{1, 2, \dots, n\}$. Here $T^c := \{1, 2, \dots, n\} - T$.

and which has basic optimal strategies that are optimal in the original game when augmented by zeros in the restored coordinates.” A *basic* optimal strategy is an extreme point of the convex set of optimal strategies. In the case of Model B3 we have not proved that our optimal strategy is basic, but a kernel nevertheless exists.

We define

$$\begin{aligned} T_0 &:= \{1 \leq l \leq n : a_{i,0}(l) < a_{i,1}(l) \text{ for } i = 0, 1, \dots, m - 1\} \\ T_1 &:= \{1 \leq l \leq n : a_{i,0}(l) > a_{i,1}(l) \text{ for } i = 0, 1, \dots, m - 1\} \\ T_* &:= \{1, 2, \dots, n\} - T_0 - T_1 \end{aligned}$$

and put $n_* := |T_*|$. Then, given $T \subset \{1, 2, \dots, n\}$, player II's pure strategy T is strictly dominated unless $T_1 \subset T \subset T_1 \cup T_*$. Therefore, the $m \times 2^n$ matrix game can be reduced to an $m \times 2^{n_*}$ matrix game.

Remark. With Z_i defined as in the first paragraph of this section, the lemma implies that, regardless of player I's strategy choice, it is optimal for player II to choose move 1 if $Z_i \in T_1$ and move 0 if $Z_i \in T_0$.

Proof. See [12] for the proof. □

We now formalize Foster's [9] algorithm for solving the 2×2^n matrix game described in the special case of Lemma 1 in which $m = 2$. (See also Foster's discussion of [15].) Of course, it is well known that a 2×2^n matrix game can be solved by maximizing the lower envelope of 2^n linear functions on $[0, 1]$. Foster's algorithm, when it applies, is much more efficient.

Lemma 2. Let $n \geq 1$ and consider a 2×2^n matrix game of the following form. Player I has two pure strategies, labeled 0 and 1. Player II has 2^n pure strategies, labeled by the subsets $T \subset \{1, 2, \dots, n\}$. For $i = 0, 1$, there exist probabilities $p_i(0) \geq 0, p_i(1) > 0, \dots, p_i(n) > 0$ with $p_i(0) + p_i(1) + \dots + p_i(n) = 1$ together with a real number $a_i(0)$, and for $l = 1, 2, \dots, n$, there exists a real 2×2 matrix

$$\begin{pmatrix} a_{0,0}(l) & a_{0,1}(l) \\ a_{1,0}(l) & a_{1,1}(l) \end{pmatrix}$$

The 2×2^n matrix game has payoff matrix with (i, T) entry given by

$$a_{i,T} := p_i(0)a_i(0) + \sum_{l \in T} p_i(l)a_{i,1}(l) + \sum_{l \in T^c} p_i(l)a_{i,0}(l)$$

for $i \in \{0, 1\}$ and $T \subset \{1, 2, \dots, n\}$.

We define

$$\begin{aligned} T_{00} &:= \{1 \leq l \leq n : a_{0,0}(l) < a_{0,1}(l) \text{ and } a_{1,0}(l) < a_{1,1}(l)\} \\ T_{01} &:= \{1 \leq l \leq n : a_{0,0}(l) < a_{0,1}(l) \text{ and } a_{1,0}(l) > a_{1,1}(l)\} \\ T_{10} &:= \{1 \leq l \leq n : a_{0,0}(l) > a_{0,1}(l) \text{ and } a_{1,0}(l) < a_{1,1}(l)\} \\ T_{11} &:= \{1 \leq l \leq n : a_{0,0}(l) > a_{0,1}(l) \text{ and } a_{1,0}(l) > a_{1,1}(l)\} \end{aligned}$$

and assume that $T_{00} \cup T_{01} \cup T_{10} \cup T_{11} = \{1, 2, \dots, n\}$. If player I uses the mixed strategy $(1 - p, p)$, then player II's best response is

$$T(p) := T_{11} \cup \{l \in T_{01} : p(l) < p\} \cup \{l \in T_{10} : p(l) > p\} \tag{2}$$

where

$$p(l) := \frac{p_0(l)[a_{0,0}(l) - a_{0,1}(l)]}{p_0(l)[a_{0,0}(l) - a_{0,1}(l)] + p_1(l)[a_{1,1}(l) - a_{1,0}(l)]}$$

This player I mixed strategy and player II best response leads to a player I expected gain of

$$\begin{aligned}
 V(p) := & (1 - p)e_0 + p e_1 + \sum_{\{l \in T_{01}: p(l) < p\} \cup \{l \in T_{10}: p(l) > p\}} [(1 - p)p_0(l)a_{0,1}(l) + p p_1(l)a_{1,1}(l)] \\
 & + \sum_{\{l \in T_{01}: p(l) \geq p\} \cup \{l \in T_{10}: p(l) \leq p\}} [(1 - p)p_0(l)a_{0,0}(l) + p p_1(l)a_{1,0}(l)] \tag{3}
 \end{aligned}$$

where

$$e_i := p_i(0)a_i(0) + \sum_{l \in T_{11}} p_i(l)a_{i,1}(l) + \sum_{l \in T_{00}} p_i(l)a_{i,0}(l)$$

The function $p \mapsto V(p)$ is the lower envelope of the family of linear functions $p \mapsto (1 - p)a_{0,T} + p a_{1,T}$, where T ranges over $T_{11} \subset T \subset T_{11} \cup T_{01} \cup T_{10}$. Therefore, the value of the game is

$$V := \max_{0 \leq p \leq 1} V(p) = \max \left(V(0), V(1), \max_{l \in T_{01} \cup T_{10}} V(p(l)) \right) = V(p^*)$$

If the last equality uniquely determines p^* and if $p^* = p(l^*)$ for a unique $l^* \in T_{01} \cup T_{10}$, then player I's unique optimal strategy is $(1 - p^*, p^*)$ and the two columns of the kernel are uniquely specified as $T(p^*)$ and $T(p^*) \cup \{l^*\}$. Their unique optimal mixture $(1 - q^*, q^*)$ is obtained by solving the 2×2 kernel.

Proof. Notice that l belongs to $T(p)$ if $l \in T_{11}$ or if both $l \in T_{01} \cup T_{10}$ and

$$(1 - p)p_0(l)a_{0,1}(l) + p p_1(l)a_{1,1}(l) < (1 - p)p_0(l)a_{0,0}(l) + p p_1(l)a_{1,0}(l)$$

implying Equation (2). The function in Equation (3) is continuous and piecewise linear, hence is maximized at 0, 1, or one of the points $p(l)$ at which its slope changes. The remaining conclusions of the lemma follow easily. □

3. Evaluation of the Payoff Matrix

We begin by considering the game under Model B3. Let $X_1 \leq X_2$ be the values of the two cards dealt to Player and $Y_1 \leq Y_2$ the values of the two cards dealt to Banker. Define the function $M : \{0, 1, \dots\} \mapsto \{0, 1, \dots, 9\}$ by $M(i) := \text{Mod}(i, 10)$. Then $X := M(X_1 + X_2)$ is Player's two-card total and $Y := M(Y_1 + Y_2)$ is Banker's two-card total. On the event $\{X \leq 7, Y \leq 7\}$, let X_3 denote the value of Player's third card if he draws, and let $X_3 := \emptyset$ if he stands. Similarly, let Y_3 denote the value of Banker's third card if he draws, and let $Y_3 := \emptyset$ if he stands.

As the rules specify, Player's pure strategies can be indexed by the sets S satisfying

$$\begin{aligned}
 & \{(i_1, i_2) : 0 \leq i_1 \leq i_2 \leq 9, M(i_1 + i_2) \leq 4\} \\
 & \subset S \subset \{(i_1, i_2) : 0 \leq i_1 \leq i_2 \leq 9, M(i_1 + i_2) \leq 5\} \tag{4}
 \end{aligned}$$

Assuming $X \leq 7$ and $Y \leq 7$, Player draws if $(X_1, X_2) \in S$ and stands otherwise. Since the set of pairs (i_1, i_2) satisfying $0 \leq i_1 \leq i_2 \leq 9$ and $M(i_1 + i_2) = 5$ contains $(0, 5), (1, 4), (2, 3), (6, 9)$, and $(7, 8)$, it follows that Player has 2^5 pure strategies.

On the other hand, Banker's pure strategies can be indexed by the sets T satisfying

$$T \subset \{(j_1, j_2) : 0 \leq j_1 \leq j_2 \leq 9, M(j_1 + j_2) \leq 7\} \times \{0, 1, 2, \dots, 9, \emptyset\} \tag{5}$$

Assuming $X \leq 7$ and $Y \leq 7$, Banker draws if $(Y_1, Y_2, X_3) \in T$ and stands otherwise. Since there are 44 pairs (j_1, j_2) satisfying $0 \leq j_1 \leq j_2 \leq 9$ and $M(j_1 + j_2) \leq 7$, and since $44 \times 11 = 484$, it follows that Banker has 2^{484} pure strategies.

Thus, *baccara* is a $2^5 \times 2^{484}$ matrix game. Let us denote by $G_{S,T}$ Player's profit from a one-unit bet when he adopts pure strategy S and Banker adopts pure strategy T , so that $a_{S,T} := E[G_{S,T}]$ is the (S, T) entry in the payoff matrix. Then

$$\begin{aligned}
 a_{S,T} &= E[G_{S,T}] = P(X \in \{8, 9\}, X > Y) - P(Y \in \{8, 9\}, Y > X) + E[G_{S,T} 1_{\{X \leq 7, Y \leq 7\}}] \\
 &= E[G_{S,T} 1_{\{X \leq 7, Y \leq 7\}}] \\
 &= \sum_{M(j_1+j_2) \leq 7} \sum_{k=0}^9 P((X_1, X_2) \in S, (Y_1, Y_2) = (j_1, j_2), X_3 = k) \\
 &\quad \cdot E[G_{S,T} \mid (X_1, X_2) \in S, (Y_1, Y_2) = (j_1, j_2), X_3 = k] \\
 &+ \sum_{M(j_1+j_2) \leq 7} P((X_1, X_2) \in S^c, (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset) \\
 &\quad \cdot E[G_{S,T} \mid (X_1, X_2) \in S^c, (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset] \tag{6}
 \end{aligned}$$

where $S^c := \{(i_1, i_2) : 0 \leq i_1 \leq i_2 \leq 9, M(i_1 + i_2) \leq 7\} - S$.

Let us now define, for S and T , for (j_1, j_2) satisfying $0 \leq j_1 \leq j_2 \leq 9$ and $M(j_1 + j_2) \leq 7$, and for $k \in \{0, 1, \dots, 9\}$,

$$\begin{aligned}
 a_{S,l}(j_1, j_2, k) &:= E[G_{S,T} \mid (X_1, X_2) \in S, (Y_1, Y_2) = (j_1, j_2), X_3 = k] \\
 a_{S,l}(j_1, j_2, \emptyset) &:= E[G_{S,T} \mid (X_1, X_2) \in S^c, (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset] \tag{7}
 \end{aligned}$$

where $l = 1$ if (j_1, j_2, k) (respectively, (j_1, j_2, \emptyset)) belongs to T ; and $l = 0$ if (j_1, j_2, k) (respectively, (j_1, j_2, \emptyset)) belongs to T^c (the complement of T relative to the set in Equation (5)). Defining also

$$\begin{aligned}
 p_S(j_1, j_2, k) &:= P((X_1, X_2) \in S, (Y_1, Y_2) = (j_1, j_2), X_3 = k) \\
 p_S(j_1, j_2, \emptyset) &:= P((X_1, X_2) \in S^c, (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset)
 \end{aligned}$$

we have, from Equation (6),

$$\begin{aligned}
 a_{S,T} &= \sum_{(j_1, j_2, k) \in T \text{ with } k \neq \emptyset} p_S(j_1, j_2, k) a_{S,1}(j_1, j_2, k) + \sum_{(j_1, j_2, k) \in T^c \text{ with } k \neq \emptyset} p_S(j_1, j_2, k) a_{S,0}(j_1, j_2, k) \\
 &+ \sum_{(j_1, j_2, \emptyset) \in T} p_S(j_1, j_2, \emptyset) a_{S,1}(j_1, j_2, \emptyset) + \sum_{(j_1, j_2, \emptyset) \in T^c} p_S(j_1, j_2, \emptyset) a_{S,0}(j_1, j_2, \emptyset) \tag{8}
 \end{aligned}$$

To evaluate the conditional expectations in Equation (7), we condition on (X_1, X_2) :

$$\begin{aligned}
 a_{S,l}(j_1, j_2, k) &= \sum_{(i_1, i_2) \in S} P((X_1, X_2) = (i_1, i_2) \mid (X_1, X_2) \in S, (Y_1, Y_2) = (j_1, j_2), X_3 = k) \\
 &\quad \cdot E[G_{S,T} \mid (X_1, X_2) = (i_1, i_2), (Y_1, Y_2) = (j_1, j_2), X_3 = k] \tag{9}
 \end{aligned}$$

if $k \neq \emptyset$ and

$$\begin{aligned}
 a_{S,l}(j_1, j_2, \emptyset) &= \sum_{(i_1, i_2) \in S^c} P((X_1, X_2) = (i_1, i_2) \mid (X_1, X_2) \in S^c, (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset) \\
 &\quad \cdot E[G_{S,T} \mid (X_1, X_2) = (i_1, i_2), (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset] \tag{10}
 \end{aligned}$$

To evaluate the conditional expectations in Equations (9) and (10), there are four cases to consider:

Case 1. $(i_1, i_2) \in S, (j_1, j_2, k) \in T$ with $k \neq \emptyset$ (both Player and Banker draw). Here, for d decks,

$$E[G_{S,T} \mid (X_1, X_2) = (i_1, i_2), (Y_1, Y_2) = (j_1, j_2), X_3 = k] = \sum_{l=0}^9 \frac{4d(1 + 3\delta_{l,0}) - \delta_{l,i_1} - \delta_{l,i_2} - \delta_{l,j_1} - \delta_{l,j_2} - \delta_{l,k}}{52d - 5} \operatorname{sgn}(M(i_1 + i_2 + k) - M(j_1 + j_2 + l)) \quad (11)$$

which converges, as $d \rightarrow \infty$, to

$$\sum_{l=0}^9 \frac{1 + 3\delta_{l,0}}{13} \operatorname{sgn}(M(i_1 + i_2 + k) - M(j_1 + j_2 + l)) \quad (12)$$

Case 2. $(i_1, i_2) \in S, (j_1, j_2, k) \in T^c$ with $k \neq \emptyset$ (Player draws, Banker stands). Regardless of the number of decks,

$$E[G_{S,T} \mid (X_1, X_2) = (i_1, i_2), (Y_1, Y_2) = (j_1, j_2), X_3 = k] = \operatorname{sgn}(M(i_1 + i_2 + k) - M(j_1 + j_2))$$

Case 3. $(i_1, i_2) \in S^c, (j_1, j_2, \emptyset) \in T$ (Player stands, Banker draws). For d decks,

$$E[G_{S,T} \mid (X_1, X_2) = (i_1, i_2), (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset] = \sum_{l=0}^9 \frac{4d(1 + 3\delta_{l,0}) - \delta_{l,i_1} - \delta_{l,i_2} - \delta_{l,j_1} - \delta_{l,j_2}}{52d - 4} \operatorname{sgn}(M(i_1 + i_2) - M(j_1 + j_2 + l)) \quad (13)$$

which converges, as $d \rightarrow \infty$, to

$$\sum_{l=0}^9 \frac{1 + 3\delta_{l,0}}{13} \operatorname{sgn}(M(i_1 + i_2) - M(j_1 + j_2 + l)) \quad (14)$$

Case 4. $(i_1, i_2) \in S^c, (j_1, j_2, \emptyset) \in T^c$ (both Player and Banker stand). Regardless of the number of decks,

$$E[G_{S,T} \mid (X_1, X_2) = (i_1, i_2), (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset] = \operatorname{sgn}(M(i_1 + i_2) - M(j_1 + j_2))$$

Next, to evaluate

$$P((X_1, X_2) = (i_1, i_2) \mid (X_1, X_2) \in S, (Y_1, Y_2) = (j_1, j_2), X_3 = k)$$

we begin with a full d -deck shoe except for three cards, one j_1 , one j_2 , and one k , removed. It will comprise m_0 0s, m_1 1s, ..., and m_9 9s, where

$$m_r := 4d(1 + 3\delta_{r,0}) - \delta_{r,j_1} - \delta_{r,j_2} - \delta_{r,k}, \quad r = 0, 1, \dots, 9$$

The number of equally likely two-card hands is $\binom{52d-3}{2}$, and the number of those that belong to S is

$$m := \sum_{(u_1, u_2) \in S} \frac{m_{u_1}(m_{u_2} - \delta_{u_2, u_1})}{1 + \delta_{u_1, u_2}}$$

Then

$$\begin{aligned}
 &P((X_1, X_2) = (i_1, i_2) \mid (X_1, X_2) \in S, (Y_1, Y_2) = (j_1, j_2), X_3 = k) \\
 &= \frac{m_{i_1}(m_{i_2} - \delta_{i_2, i_1})}{1 + \delta_{i_1, i_2}} \frac{1_S((i_1, i_2))}{m}
 \end{aligned} \tag{15}$$

and

$$\begin{aligned}
 p_S(j_1, j_2, k) &= P((Y_1, Y_2) = (j_1, j_2), X_3 = k)P((X_1, X_2) \in S \mid (Y_1, Y_2) = (j_1, j_2), X_3 = k) \\
 &= \frac{(2 - \delta_{j_1, j_2})4d(1 + 3\delta_{j_1, 0})[4d(1 + 3\delta_{j_2, 0}) - \delta_{j_2, j_1}]}{(52d)(52d - 1)} \\
 &\quad \cdot \frac{4d(1 + 3\delta_{k, 0}) - \delta_{k, j_1} - \delta_{k, j_2}}{52d - 2} \frac{m}{\binom{52d-3}{2}}
 \end{aligned} \tag{16}$$

Finally, to evaluate

$$P((X_1, X_2) = (i_1, i_2) \mid (X_1, X_2) \in S^c, (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset)$$

we begin with a full d -deck shoe except for two cards, one j_1 and one j_2 , removed. It will comprise m'_0 0s, m'_1 1s, ..., and m'_9 9s, where

$$m'_r := 4d(1 + 3\delta_{r, 0}) - \delta_{r, j_1} - \delta_{r, j_2}, \quad r = 0, 1, \dots, 9$$

The number of equally likely two-card hands is $\binom{52d-2}{2}$, and the number of those that belong to S^c is

$$m' := \sum_{(u_1, u_2) \in S^c} \frac{m'_{u_1}(m'_{u_2} - \delta_{u_2, u_1})}{1 + \delta_{u_1, u_2}}$$

Then

$$\begin{aligned}
 &P((X_1, X_2) = (i_1, i_2) \mid (X_1, X_2) \in S^c, (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset) \\
 &= \frac{m'_{i_1}(m'_{i_2} - \delta_{i_2, i_1})}{1 + \delta_{i_1, i_2}} \frac{1_{S^c}((i_1, i_2))}{m'}
 \end{aligned} \tag{17}$$

and

$$\begin{aligned}
 p_S(j_1, j_2, \emptyset) &= P((Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset)P((X_1, X_2) \in S^c \mid (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset) \\
 &= \frac{(2 - \delta_{j_1, j_2})4d(1 + 3\delta_{j_1, 0})[4d(1 + 3\delta_{j_2, 0}) - \delta_{j_2, j_1}]}{(52d)(52d - 1)} \frac{m'}{\binom{52d-2}{2}}
 \end{aligned} \tag{18}$$

This suffices to complete the evaluation of Equations (9) and (10) when cards are dealt without replacement from a d -deck shoe.

The assumption that cards are dealt *with* replacement from a single deck can be modeled by letting $d \rightarrow \infty$ in the assumption that cards are dealt without replacement from a d -deck shoe. The formulas are simpler in this case:

$$\begin{aligned}
 &P((X_1, X_2) = (i_1, i_2) \mid (X_1, X_2) \in S, (Y_1, Y_2) = (j_1, j_2), X_3 = k) \\
 &= \frac{(2 - \delta_{i_1, i_2})(1 + 3\delta_{i_1, 0})(1 + 3\delta_{i_2, 0})1_S((i_1, i_2))}{89 + 8|S \cap \{(0, 5)\}| + 2|S \cap \{(1, 4), (2, 3), (6, 9), (7, 8)\}|}
 \end{aligned} \tag{19}$$

$$p_S(j_1, j_2, k) = \frac{(2 - \delta_{j_1, j_2})(1 + 3\delta_{j_1, 0})(1 + 3\delta_{j_2, 0})}{(13)^2} \frac{1 + 3\delta_{k, 0}}{13} \cdot \frac{89 + 8|S \cap \{(0, 5)\}| + 2|S \cap \{(1, 4), (2, 3), (6, 9), (7, 8)\}|}{(13)^2} \tag{20}$$

$$P((X_1, X_2) = (i_1, i_2) \mid (X_1, X_2) \in S^c, (Y_1, Y_2) = (j_1, j_2), X_3 = \emptyset) = \frac{(2 - \delta_{i_1, i_2})(1 + 3\delta_{i_1, 0})(1 + 3\delta_{i_2, 0})1_{S^c}((i_1, i_2))}{32 + 8|S^c \cap \{(0, 5)\}| + 2|S^c \cap \{(1, 4), (2, 3), (6, 9), (7, 8)\}|} \tag{21}$$

$$p_S(j_1, j_2, \emptyset) = \frac{(2 - \delta_{j_1, j_2})(1 + 3\delta_{j_1, 0})(1 + 3\delta_{j_2, 0})}{(13)^2} \cdot \frac{32 + 8|S^c \cap \{(0, 5)\}| + 2|S^c \cap \{(1, 4), (2, 3), (6, 9), (7, 8)\}|}{(13)^2} \tag{22}$$

Here 89 comes from 25 + 16 + 16 + 16 + 16, where the summands correspond to totals 0, 1, 2, 3, 4; 32 is 16 + 16, corresponding to totals 6 and 7.

In summary, we can evaluate Equation (8) under Model B3 or A3. Restricting S to the two extremes in Equation (4), we obtain Equation (8) under Model B2 or A2 as a special case. Finally, for Models B1 and A1, we can derive the analogue of Equation (8) from results already obtained. Specifically,

$$a_{S^\circ, T^\circ} = \sum_{(j, k) \in T^\circ \text{ with } k \neq \emptyset} p_{S^\circ}(j, k) a_{S^\circ, 1}(j, k) + \sum_{(j, k) \in (T^\circ)^c \text{ with } k \neq \emptyset} p_{S^\circ}(j, k) a_{S^\circ, 0}(j, k) + \sum_{(j, \emptyset) \in T^\circ} p_{S^\circ}(j, \emptyset) a_{S^\circ, 1}(j, \emptyset) + \sum_{(j, \emptyset) \in (T^\circ)^c} p_{S^\circ}(j, \emptyset) a_{S^\circ, 0}(j, \emptyset)$$

where

$$a_{S^\circ, l}(j, k) := P(G_{S^\circ, T^\circ} \mid X \in S^\circ, Y = j, X_3 = k) = \sum_{M(j_1 + j_2) = j} p_S(j_1, j_2, k) a_{S, l}(j_1, j_2, k) / p_{S^\circ}(j, k)$$

$$a_{S^\circ, l}(j, \emptyset) := P(G_{S^\circ, T^\circ} \mid X \in (S^\circ)^c, Y = j, X_3 = \emptyset) = \sum_{M(j_1 + j_2) = j} p_S(j_1, j_2, \emptyset) a_{S, l}(j_1, j_2, \emptyset) / p_{S^\circ}(j, \emptyset)$$

$$p_{S^\circ}(j, k) = P(X \in S^\circ, Y = j, X_3 = k) = \sum_{M(j_1 + j_2) = j} p_S(j_1, j_2, k)$$

and

$$p_{S^\circ}(j, \emptyset) = P(X \in (S^\circ)^c, Y = j, X_3 = \emptyset) = \sum_{M(j_1 + j_2) = j} p_S(j_1, j_2, \emptyset)$$

here $S^\circ = \{0, 1, 2, 3, 4\}$ or $\{0, 1, 2, 3, 4, 5\}$, $(S^\circ)^c = \{0, 1, \dots, 7\} - S^\circ$, and $T^\circ \subset \{0, 1, \dots, 7\} \times \{0, 1, \dots, 9, \emptyset\}$, while S and T are the corresponding subsets of $\{(i_1, i_2) : 0 \leq i_1 \leq i_2 \leq 9, M(i_1 + i_2) \leq 7\}$ and $\{(j_1, j_2) : 0 \leq j_1 \leq j_2 \leq 9, M(j_1 + j_2) \leq 7\} \times \{0, 1, \dots, 9, \emptyset\}$, respectively.

4. Banker’s Strictly Dominated Pure Strategies

Our next step is to show that Lemma 1 applies (with Player and Banker playing the roles of player I and player II, respectively), allowing us to reduce the game to a more manageable size. The payoff matrix in Equation (8) has the form in Equation (1) with $m = 32$, $n = 484$, $p_i(0) = P(X \in \{8, 9\} \text{ or } Y \in \{8, 9\})$, and $a_i(0) = 0$. It remains to evaluate T_0 , T_1 , and T_* of the lemma.

Results are summarized in Table 1. T_1 (respectively, T_0) is the set of triples (j_1, j_2, k) for which $a_{S,1}(j_1, j_2, k) < a_{S,0}(j_1, j_2, k)$ (respectively, $>$) for each of Player’s 2^5 pure strategies S , indicated by a D (respectively, S) in the corresponding entry of the table. T_* is the remaining set of triples (j_1, j_2, k) , indicated by a * in the corresponding entry of the table. Of particular interest is $n_d := |T_*|$.

Theorem 3. (a) Under Model B3 with the number of decks being a positive integer d , Lemma 1 applies. The sets T_0 , T_1 , and T_* of the lemma can be inferred from Table 1, with entries S, D, and * located at elements of T_0 , T_1 , and T_* , respectively. In particular, for $d = 1, 2, \dots, 10$, $n_d = 23, 21, 20, 19, 19, 18, 21, 23, 23, 23$, respectively, and, if $d \geq 11$, $n_d = 22$.

(b) Exactly the same conclusions hold under Model B2.

(c) Lemma 1 applies under Models A2 and A3, with the same results as those under Models B2 and B3 with $d \geq 11$.

Proof. It will occasionally be convenient to label the 2^5 choices of S by the integers 0 to 31. Strategy 0 (respectively, strategy 31) denotes Player’s pure strategy of standing (respectively, drawing) on a two-card total of 5, regardless of its composition. More generally, strategy $u \in \{0, 1, \dots, 31\}$ is specified by the 5-bit binary representation of u . For example, strategy 19 (binary 10011) corresponds to drawing on (0, 5), standing on (1, 4) and (2, 3), and drawing on (6, 9) and (7, 8).

Table 1 is identical under Models B2 and B3 because, defining

$$b_u(j_1, j_2, k) := a_{u,1}(j_1, j_2, k) - a_{u,0}(j_1, j_2, k)$$

for $u = 0, 1, \dots, 31$, $0 \leq j_1 \leq j_2 \leq 9$ with $M(j_1 + j_2) \leq 7$, and $k = 0, 1, \dots, 9, \emptyset$, we have, with a few exceptions,

$$b_u(j_1, j_2, k) \in [\min\{b_0(j_1, j_2, k), b_{31}(j_1, j_2, k)\}, \max\{b_0(j_1, j_2, k), b_{31}(j_1, j_2, k)\}]$$

for $u = 1, 2, \dots, 30$. The exceptions occur only when $d = 1$ and only when $(j_1, j_2, k) = (0, 0, 9)$, $(5, 5, 9)$, or $(5, 6, 0)$. (See Section 7 for an explanation of why this is to be expected.)

It follows that, if $b_0(j_1, j_2, k) < 0$ and $b_{31}(j_1, j_2, k) < 0$, then the (j_1, j_2, k) entry in Table 1 is D (draw); in 301 of the 484 entries, this property holds for every $d \geq 1$. If $b_0(j_1, j_2, k) > 0$ and $b_{31}(j_1, j_2, k) > 0$, then the (j_1, j_2, k) entry in Table 1 is S (stand); in 151 of the 484 entries, this property holds for every $d \geq 1$. If $b_0(j_1, j_2, k) > 0 > b_{31}(j_1, j_2, k)$ or $b_0(j_1, j_2, k) < 0 < b_{31}(j_1, j_2, k)$, then the (j_1, j_2, k) entry in Table 1 is *; in 13 of the 484 entries, this property holds for every $d \geq 1$. This accounts for all but 19 entries in Table 1, those marked with footnotes, in which the sign of $b_0(j_1, j_2, k)$ or $b_{31}(j_1, j_2, k)$ depends on d .

For example, one of the 19 is (3, 3, 6). Indeed,

$$b_{31}(3, 3, 6) = -\frac{2(80d^3 - 832d^2 + 135d - 2)}{(52d - 5)(840d^2 - 114d + 1)} > 0 \text{ (respectively, } < 0)$$

for all $d \leq 10$ (respectively, $d \geq 11$) and

$$b_0(3, 3, 6) = -\frac{2(848d^3 - 952d^2 + 135d - 2)}{(52d - 5)(712d^2 - 102d + 1)} < 0$$

for all $d \geq 1$. It follows that the $(3, 3, 6)$ entry of Table 1 is * if $d \leq 10$ and D if $d \geq 11$. The other 483 cases are analyzed similarly. □

Table 1. Banker’s optimal move (preliminary version) under Model B2 or B3 with $d = 6$, indicated by D (draw) or S (stand), except in the $n_6 = 18$ cases indicated by * in which it depends on Player’s strategy. Adjustments to the table for other positive integers d are specified by footnotes. Under Model A2 or A3, results are the same as those under Model B2 or B3 with $d \geq 11$.

Banker’s Two-Card		Player’s Third Card (\emptyset if Player stands)										
Total	Hand	0	1	2	3	4	5	6	7	8	9	\emptyset
0, 1, 2		D	D	D	D	D	D	D	D	D	D	D
3	(0, 3)	D	D	D	D	D	D	D	D	S ²	*	D
3	(1, 2)	D	D	D	D	D	D	D	D	S ¹	*	D
3	(4, 9)	D	D	D	D	D	D	D	D	S ⁵	*	D
3	(5, 8)	D	D	D	D	D	D	D	D	S ³	*	D
3	(6, 7)	D	D	D	D	D	D	D	D	S ³	*	D
4	(0, 4)	S	S ⁸	D ¹	D	D	D	D	D	S	S	D
4	(1, 3)	S	S ⁷	D ¹	D	D	D	D	D	S	S	D
4	(2, 2)	S	* ³	D ¹	D	D	D	D	D	S	S	D
4	(5, 9)	S	S ⁷	D ¹	D	D	D	D	D	S	S	D
4	(6, 8)	S	*	D	D	D	D	D	D	S	S	D
4	(7, 7)	S	*	D	D	D	D	D	D	S	S	D
5	(0, 5)	S	S	S	S	* ¹	D	D	D	S	S	D
5	(1, 4)	S	S	S	S	S ⁷	D	D	D	S	S	D
5	(2, 3)	S	S	S	S	S ⁸	D	D	D	S	S	D
5	(6, 9)	S	S	S	S	* ¹	D	D	D	S	S	D
5	(7, 8)	S	S	S	S	* ¹	D	D	D	S	S	D
6	(0, 6)	S	S	S	S	S	S	D	D	S	S	*
6	(1, 5)	S	S	S	S	S	S	D	D	S	S	*
6	(2, 4)	S	S	S	S	S	S	D	D	S	S	*
6	(3, 3)	S	S	S	S	S	S	* ¹¹	D	S	S	*
6	(7, 9)	S	S	S	S	S	S	D	D	S	S	*
6	(8, 8)	S	S	S	S	S	S	D	D	S	S	*
7		S	S	S	S	S	S	S	S	S	S	S

¹ Replace S by * and D by * and * by S if $d = 1$; ² Replace S by * if $d \leq 2$; ³ Replace S by * and * by S if $d \leq 3$; ⁵ Replace S by * if $d \leq 5$; ⁷ Replace S by * if $d \geq 7$; ⁸ Replace S by * if $d \geq 8$; ¹¹ Replace * by D if $d \geq 11$.

Requiring that Banker make the optimal move in each of the cases that do not depend on Player’s strategy, we have reduced the game, under Model B3 (respectively, B2), to a $2^5 \times 2^{n_d}$ (respectively, 2×2^{n_d}) matrix game, where $18 \leq n_d \leq 23$.

We have similar results for Model B1. Again, $n_d := |T_*|$.

Theorem 4. (a) Under Model B1 with the number of decks being a positive integer d , Lemma 1 applies. The sets T_0 , T_1 , and T_* of the lemma can be inferred from Table 2, with entries S, D, and * located at elements of T_0 , T_1 , and T_* , respectively. In particular, $n_1 = 4$, $n_2 = 3$, and $n_d = 4$ for all $d \geq 3$. (b) Lemma 1 applies under Model A1, with the same results as those under Model B1 with $d \geq 4$.

Table 2. Banker’s optimal move (preliminary version) under Model B1 with $d = 6$, indicated by D (draw) or S (stand), except in the $n_6 = 4$ cases indicated by * in which it depends on Player’s strategy. Adjustments to the table for other positive integers d are specified by footnotes. Under Model A1, results are the same as those under Model B1 with $d \geq 4$.

Banker’s Two-Card	Player’s Third Card (\emptyset if Player Stands)											
	0	1	2	3	4	5	6	7	8	9	\emptyset	
0, 1, 2	D	D	D	D	D	D	D	D	D	D	D	D
3	D	D	D	D	D	D	D	D	S ³	*	D	D
4	S	* ³	D ¹	D	D	D	D	D	S	S	D	D
5	S	S	S	S	* ²	D	D	D	S	S	D	D
6	S	S	S	S	S	S	D	D	S	S	*	D
7	S	S	S	S	S	S	S	S	S	S	S	S

¹ Replace D by * if $d = 1$; ² Replace * by S if $d \leq 2$; ³ Replace S by * and * by S if $d \leq 3$.

Proof. Let

$$b_u(j, k) := a_{u,1}(j, k) - a_{u,0}(j, k)$$

where $u = 0$ corresponds to $S^\circ = \{0, 1, 2, 3, 4\}$ and $u = 1$ corresponds to $S^\circ = \{0, 1, 2, 3, 4, 5\}$. If $b_u(j, k) < 0$ for $u = 0, 1$, then the (j, k) entry in Table 2 is D; in 54 of the 88 entries, this property holds for every $d \geq 1$. If $b_u(j, k) > 0$ for $u = 0, 1$, then the (j, k) entry in Table 2 is S; in 28 of the 88 entries, this property holds for every $d \geq 1$. If $b_0(j, k) > 0 > b_1(j, k)$ or $b_0(j, k) < 0 < b_1(j, k)$, then the (j, k) entry in Table 1 is *; in two of the 88 entries, namely (3, 9) and (6, \emptyset), this property holds for every $d \geq 1$. This accounts for all but four entries in Table 2, those marked with footnotes, in which the sign of $b_u(j, k)$ depends on d for $u = 0$ or $u = 1$.

For example, one of the four is (5, 4). Indeed,

$$b_1(5, 4) = -\frac{15\,360\,d^4 - 45\,184\,d^3 + 9040\,d^2 - 588\,d + 13}{(52\,d - 5)(26\,880\,d^3 - 4680\,d^2 + 242\,d - 3)} < 0 \text{ (respectively, } > 0)$$

for $d \geq 3$ (respectively, $d \leq 2$), and

$$b_0(5, 4) = \frac{1024\,d^4 + 37\,248\,d^3 - 7792\,d^2 + 492\,d - 7}{(52\,d - 5)(22\,784\,d^3 - 3976\,d^2 + 194\,d - 1)} > 0$$

for all $d \geq 1$. Therefore, the entry in the (5, 4) position of Table 2 is * if $d \geq 3$ and S if $d \leq 2$. The other 87 cases are analyzed similarly. \square

5. Solutions under Models A1, A2, and A3

We recall Kemeny and Snell’s [7] solution of the 2×2^{88} matrix game that assumes Model A1. (See Deloche and Oguer 2007 for an alternative approach based on the extensive, rather than the strategic, form of the game.) Implicitly using Lemma 1, they reduced the number of Banker pure strategies from 2^{88} to just 2^4 (see Theorem 4). The 2×2 kernel of the game was determined to be

$$\begin{matrix} & \text{B: S on 6, } \emptyset & \text{B: D on 6, } \emptyset \\ \text{P: S on 5} & \begin{pmatrix} -4564 & -2692 \\ -3705 & -4121 \end{pmatrix} & \\ \text{P: D on 5} & & \end{matrix} 2^4 / (13)^6 \tag{23}$$

implying that the following Player and Banker strategies are uniquely optimal. Player draws on a two-card total of 5 with probability

$$p = \frac{9}{11} \approx 0.818182 \tag{24}$$

and Banker draws on a two-card total of 6, when Player stands, with probability

$$q = \frac{859}{2288} \approx 0.375437 \tag{25}$$

The value of the game (to Player) is

$$v = -\frac{679\,568}{53\,094\,899} \approx -0.0127991 \tag{26}$$

The fully specified optimal strategy for Banker is given in Table 3 with $d \geq 4$ and q as in Equation (25).

Table 3. Banker’s optimal move (final version) under Model B1 with d being a positive integer, indicated by D (draw) or S (stand), or (S, D) (stand with probability $1 - q$, draw with probability q). Here q is as in Equations (33) and (34). Under Model A1, results are the same as those under Model B1 with $d \geq 4$, except that q is as in Equation (25).

Banker’s Two-Card Total	Player’s Third Card (\emptyset if Player Stands)											
	0	1	2	3	4	5	6	7	8	9	\emptyset	
0, 1, 2	D	D	D	D	D	D	D	D	D	D	D	D
3	D	D	D	D	D	D	D	D	S	D	D	
4	S	S	D	D	D	D	D	D	S	S	D	
5	S	S	S	S	*	D	D	D	S	S	D	
6	S	S	S	S	S	S	D	D	S	S	(S, D)	
7	S	S	S	S	S	S	S	S	S	S	S	

* D if $d \geq 4$, S if $d \leq 3$.

Let us extend this analysis from Model A1 to Model A3. Again we have a $2^5 \times 2^{484}$ matrix game, and the payoff matrix can be evaluated as in Section 4, using Equations (12), (14), and (19)–(22) in place of

Equations (11), (13), and (15)–(18). We can apply Lemma 1 and reduce the game to a $2^5 \times 2^{22}$ matrix game. We obtain the special case of Table 1 in which $d \geq 11$.

When we evaluate this $2^5 \times 2^{22}$ payoff matrix, we find that a number of rows are identical. When several rows are identical, we eliminate duplicates. When we make this reduction and rearrange the remaining rows in a more natural order, we are left with 9 rows, labeled by 0–8, which have a special structure. Specifically, row i corresponds to Player’s mixed strategy (under Model A1) of drawing on a two-card total of 5 with probability $i/8$. The reason that multiples of $1/8$ appear is that, given that Player has a two-card total of 5, he has (0, 5), (1, 4), (2, 3), (6, 9), or (7, 8) with probabilities $4/8, 1/8, 1/8, 1/8,$ and $1/8$, respectively.

There are also a number of identical columns. When we apply a similar reduction and rearrangement to the columns, we are left with $9^2(17)^2 = 23\,409$ columns, labeled by

$$\{0, 1, \dots, 8\} \times \{0, 1, \dots, 16\} \times \{0, 1, \dots, 8\} \times \{0, 1, \dots, 16\}$$

which have a similar structure. Specifically, column (j_1, j_2, j_3, j_4) corresponds to Banker’s mixed strategy (under Model A1) of drawing on a two-card total of 3, when Player’s third card is 9, with probability $j_1/8$; of drawing on a two-card total of 4, when Player’s third card is 1, with probability $j_2/16$; of drawing on a two-card total of 5, when Player’s third card is 4, with probability $j_3/8$; and of drawing on a two-card total of 6, when Player stands, with probability $j_4/16$. The reason that multiples of $1/8$ or $1/16$ appear is that, given that Banker has a two-card total of 3, he has (0, 3), (1, 2), (4, 9), (5, 8), or (6, 7) with probabilities $4/8, 1/8, 1/8, 1/8,$ and $1/8$, respectively; given that Banker has a two-card total of 4, he has (0, 4), (1, 3), (2, 2), (5, 9), (6, 8), or (7, 7) with probabilities $8/16, 2/16, 1/16, 2/16, 2/16,$ and $1/16$, respectively; given that Banker has a two-card total of 5, he has (0, 5), (1, 4), (2, 3), (6, 9), or (7, 8) with probabilities $4/8, 1/8, 1/8, 1/8,$ and $1/8$, respectively; given that Banker has a two-card total of 6, he has (0, 6), (1, 5), (2, 4), (3, 3), (7, 9), or (8, 8) with probabilities $8/16, 2/16, 2/16, 1/16, 2/16,$ and $1/16$, respectively.

Next, we observe that column (j_1, j_2, j_3, j_4) is a mixture of the 2^4 (Model A1) pure strategies of Banker that remain after application of Lemma 1. By the results for Model A1, optimal strategies for Banker must satisfy $j_1 = 8, j_2 = 0,$ and $j_3 = 8$. This reduces the game to a 9×17 matrix game, whose columns we relabel as 0–16. Specifically, column j corresponds to Banker’s mixed strategy (under Model A1) of drawing on a two-card total of 6, when Player stands, with probability $j/16$.

Finally, what is the solution of the 9×17 game? We have seen that rows 1–7 (respectively, columns 1–15) are mixtures of rows 0 and 8 (respectively, columns 0 and 16). In particular, rows 1–7 and columns 1–15 are dominated, but not strictly dominated. Eliminating these rows and columns results in a 2×2 matrix game, namely the kernel in Equation (23). But eliminating dominated, but not strictly dominated, rows and columns may result in a loss of solutions, and it does so in this case. Indeed, there are many solutions. For Player, given two pure strategies $i, i' \in \{0, 1, 2, \dots, 8\}$ with

$$\frac{i}{8} < \frac{9}{11} < \frac{i'}{8} \tag{27}$$

there is a unique $p \in (0, 1)$ such that

$$(1 - p) \left(\frac{i}{8} \right) + p \left(\frac{i'}{8} \right) = \frac{9}{11}$$

and the $(1 - p, p)$ -mixture of pure strategies i and i' is optimal for Player. There are 7 choices of i and 2 choices of i' that satisfy Equation (27), hence 14 pairs (i, i') that meet this condition.

For Banker, given two pure strategies $j, j' \in \{0, 1, 2, \dots, 16\}$ with

$$\frac{j}{16} < \frac{859}{2288} < \frac{j'}{16} \tag{28}$$

there is a unique $q \in (0, 1)$ such that

$$(1 - q)\left(\frac{j}{16}\right) + q\left(\frac{j'}{16}\right) = \frac{859}{2288}$$

and the $(1 - q, q)$ -mixture of pure strategies j and j' is optimal for Banker. There are 7 choices of j and 10 choices of j' that satisfy Equation (28), hence 70 pairs (j, j') that meet this condition.

Each such pair (i, i') can be combined with each such pair (j, j') , so there are $14 \times 70 = 980$ pairs of optimal strategies of this form. These are the extreme points of the convex set of equilibria. All 980 of them appear when Savani’s [16] game solver is applied.

Let us single out one of them. Take $i = 6$ and $i' = 7$, getting $p = 6/11$, and take $j = 1$ and $j' = 9$, getting $q = 179/286$. This does not uniquely determine a pair of optimal strategies because of the duplicate rows and columns that were eliminated, but one pair of optimal mixed strategies to which this corresponds is shown in Table 4.

Table 4. A pair of optimal mixed strategies in Model A3. For Banker’s fully specified optimal strategy, see Table 5 with $d \geq 10$.

Player’s Two-Card Total is 5	
(0, 5), (6, 9), (7, 8)	D
(1, 4)	(S, D) with (5/11, 6/11)
(2, 3)	S
Banker’s Two-Card Total is 6 and Player Stands	
(0, 6)	(S, D) with (107/286, 179/286)
(1, 5), (2, 4), (3, 3), (7, 9)	S
(8, 8)	D

As we will see, this pair of optimal mixed strategies is the limiting pair of optimal mixed strategies under Model B3 as $d \rightarrow \infty$.

Foster [9] remarked, “It is an interesting fact that this [optimal Player mixed] strategy is often attained approximately in practice by standing on the pair 2, 3 and calling [*i.e.*, drawing] on any other combination adding to 5; this gives approximately the right frequency of calling [*i.e.*, drawing].” Actually, it gives a drawing probability of 7/8, not the required 9/11. But, as Table 4 suggests, Player should stand also on (1, 4) with probability 5/11. Then the probability of Player drawing on a two-card total of 5 is

$$\frac{3}{4} + \frac{1}{8} \frac{6}{11} = \frac{9}{11}$$

A similar analysis applies under Model A2. Lemma 1 reduces the 2×2^{484} matrix game to 2×2^{22} . Eliminating duplicate columns reduces the game to $2 \times 9^2(17)^2$, and finally using the optimal solution under Model A1, we are left with a 2×17 matrix game. This is the game identified in [10]. As above, there are 70 extremal solutions, and they all appear when Savani’s [16] game solver is applied.

Again, we single out one of them. Player draws on a two-card total of 5 with probability 9/11, and Banker follows Table 4. As we will see, this pair of optimal mixed strategies is the limiting pair of optimal mixed strategies under Model B2 as $d \rightarrow \infty$.

6. Solution under Model B1

Recall that Table 2 applies under Model B1. See Theorem 4.

In the case $d = 6$, the 2×2^4 matrix game has kernel given by

$$\begin{matrix} & \text{B: S on 6, } \emptyset & \text{B: D on 6, } \emptyset \\ \text{P: S on 5} & \left(\begin{matrix} -23\ 256\ 431\ 632 & -13\ 884\ 629\ 124 \\ -18\ 880\ 657\ 128 & -21\ 061\ 456\ 188 \end{matrix} \right) & / 1\ 525\ 814\ 595\ 305 \\ \text{P: D on 5} & & \end{matrix}$$

implying that the following Player and Banker strategies are uniquely optimal. Player draws on a two-card total of 5 with probability

$$p_6 = \frac{7\ 631\ 761}{9\ 407\ 656} \approx 0.811229 \tag{29}$$

and Banker draws on a two-card total of 6, when Player stands, with probability

$$q_6 = \frac{546\ 971\ 813}{1\ 444\ 075\ 196} \approx 0.378770 \tag{30}$$

The value of the game (to Player) is

$$v_6 = -\frac{23\ 174\ 205\ 422\ 119\ 131}{1\ 794\ 292\ 354\ 051\ 081\ 885} \approx -0.0129155 \tag{31}$$

The fully specified optimal strategy for Banker is given in Table 3. Comparing the solution under Model A1 with that under Model B1 reveals the effect of the “with replacement” assumption. The solutions are identical except for the three parameters [Equations (24)–(26) vs. Equations (29)–(31)]. For each parameter, the relative error is less than one percent.

This analysis extends to every positive integer d .

Theorem 5. *Under Model B1 with d being a positive integer, the following Player and Banker strategies are uniquely optimal. Player draws on a two-card total of 5 with probability*

$$p_d = \frac{(36\ 864\ d^3 - 9312\ d^2 + 732\ d - 23)}{8(5632\ d^3 - 1138\ d^2 + 69\ d - 1)}, \quad d \geq 1 \tag{32}$$

and Banker draws on on a two-card total of 6, when Player stands, with probability

$$q_d = \frac{224\ 000\ d^4 - 55\ 712\ d^3 + 2936\ d^2 + 163\ d - 14}{2(52\ d - 5)(5632\ d^3 - 1138\ d^2 + 69\ d - 1)}, \quad 1 \leq d \leq 3 \tag{33}$$

or

$$q_d = \frac{439\,808\,d^4 - 107\,456\,d^3 + 5248\,d^2 + 374\,d - 31}{4(52\,d - 5)(5632\,d^3 - 1138\,d^2 + 69\,d - 1)}, \quad d \geq 4 \tag{34}$$

The value of the game (to Player) is

$$v_d = -32\,d^2(44\,396\,707\,840\,d^7 - 18\,908\,426\,240\,d^6 + 3\,279\,293\,696\,d^5 - 294\,129\,728\,d^4 + 14\,418\,160\,d^3 - 407\,352\,d^2 + 9543\,d - 220)/[(5632\,d^3 - 1138\,d^2 + 69\,d - 1)(52\,d)_6], \quad 1 \leq d \leq 3$$

or

$$v_d = -16\,d^2(89\,072\,336\,896\,d^7 - 38\,873\,874\,432\,d^6 + 6\,969\,345\,536\,d^5 - 655\,761\,920\,d^4 + 34\,638\,784\,d^3 - 1\,090\,952\,d^2 + 26\,286\,d - 537)/[(5632\,d^3 - 1138\,d^2 + 69\,d - 1)(52\,d)_6], \quad d \geq 4 \tag{35}$$

where $(52\,d)_6 := (52\,d)(52\,d - 1) \cdots (52\,d - 5)$. The fully specified optimal strategy for Banker is given in Table 3.

Proof. For $d \geq 4$, the kernel of the 2×2^4 payoff matrix is given by columns 10 and 11 (when columns are labeled from 0 to 15), namely

	B: S on 6, \emptyset	B: D on 6, \emptyset	
P: S on 5	$4(1\,168\,384\,d^4 - 284\,720\,d^3 + 22\,320\,d^2 - 446\,d - 11)$	$2\,756\,608\,d^4 - 470\,336\,d^3 + 4656\,d^2 + 3072\,d - 159$	$\frac{(-64\,d^2)}{(52\,d)_6}$
P: D on 5	$2(1\,896\,960\,d^4 - 461\,984\,d^3 + 39\,392\,d^2 - 1266\,d + 9)$	$4\,219\,904\,d^4 - 954\,112\,d^3 + 68\,384\,d^2 - 852\,d - 57$	

This implies Equation (32) for $d \geq 4$ and Equations (34) and (35).

To confirm this, we must show that, with \mathbf{A} denoting the 2×2^4 payoff matrix, we have

$$(1 - p_d, p_d)\mathbf{A} \geq (v_d, v_d, \dots, v_d)$$

This involves checking 16 inequalities (of which two are automatic). For example, the eighth and ninth components of $(1 - p_d, p_d)\mathbf{A} - (v_d, v_d, \dots, v_d)$ equal

$$16\,d^2(278\,921\,216\,d^7 - 1\,057\,021\,952\,d^6 + 410\,758\,144\,d^5 - 67\,502\,464\,d^4 + 5\,802\,464\,d^3 - 276\,248\,d^2 + 7200\,d - 97)/[(5632\,d^3 - 1138\,d^2 + 69\,d - 1)(52\,d)_6]$$

which is positive for $d \geq 4$ and negative for $1 \leq d \leq 3$. The 10th and 11th components are 0, of course. The remaining components are positive for every $d \geq 1$.

A similar analysis can then be carried out for $1 \leq d \leq 3$, in which case the kernel is given by columns 8 and 9 (of 0–15). □

We notice that the above kernel converges, as $d \rightarrow \infty$, to the kernel in Equation (23).

7. Solution under Model B2

The result that Table 1 is identical under Models B2 and B3 is less surprising than it may first appear to be. In Section 5 we saw that, under Model A3 with Player’s pure strategies labeled from 0 to 31, pure strategy $u \in \{1, 2, \dots, 30\}$ is a $(1 - p, p)$ mixture of pure strategies 0 and 31, where

$$p = \frac{4u_1 + u_2 + u_3 + u_4 + u_5}{8} \in \left\{ \frac{1}{8}, \frac{2}{8}, \dots, \frac{7}{8} \right\} \tag{36}$$

here $u_1u_2u_3u_4u_5$ is the binary form of u , that is, $u_1, u_2, u_3, u_4, u_5 \in \{0, 1\}$ and $u = 16u_1 + 8u_2 + 4u_3 + 2u_4 + u_5$. Consequently,

$$a_{u,l}(j_1, j_2, k) \text{ lies between } a_{0,l}(j_1, j_2, k) \text{ and } a_{31,l}(j_1, j_2, k) \tag{37}$$

for all $u \in \{1, 2, \dots, 30\}$, $l = 0, 1$, $0 \leq j_1 \leq j_2 \leq 9$ with $M(j_1 + j_2) \leq 7$, and $k = 0, 1, \dots, 9, \emptyset$. Under Model B3, the conditional expectations in Equation (37) should not differ much from their Model A3 counterparts, especially for large d , hence we would expect that Equation (37) holds with few exceptions. In fact, the *only* exceptions occur when $l = 1$, $M(j_1 + j_2) = 2$, $k = 8$, and $d \leq 7$ because in these cases, $a_{0,l}(j_1, j_2, k)$ and $a_{31,l}(j_1, j_2, k)$ are very close. When we consider the differences $b_u(j_1, j_2, k)$, there are even fewer exceptions, as noted previously.

One might ask why Model B2 was even considered by Downton and Lockwood [8], inasmuch as its asymmetric assumption about the available information (beyond the asymmetry inherent in the rules) may seem contrived. The answer, we believe, is that there already existed an algorithm, due to Foster [9], for solving such games. That algorithm was formalized in Lemma 2 of Section 2.

Let us recall Downton and Lockwood’s [8] solution of the 2×2^{484} matrix game that assumes Model B2. In the case $d = 6$, the 2×2 kernel of the game is found to be

$$\begin{array}{l} \text{B: S on } (0, 6), \emptyset \quad \text{B: D on } (0, 6), \emptyset \\ \text{P: S on } 5 \quad \left(\begin{array}{cc} -22\,721\,165\,499 & -18\,033\,241\,115 \\ -19\,018\,265\,931 & -20\,151\,297\,323 \end{array} \right) / 1\,525\,814\,595\,305 \\ \text{P: D on } 5 \end{array}$$

implying that the following Player and Banker strategies are uniquely optimal. Player draws on a two-card total of 5 with probability

$$p_6 = \frac{477\,191}{592\,524} \approx 0.805353$$

and Banker draws on $(0, 6)$, when Player stands, with probability

$$q_6 = \frac{77\,143\,741}{121\,269\,912} \approx 0.636133$$

The value of the game (to Player) is

$$v_6 = -\frac{974\,653\,793\,197\,999}{75\,340\,147\,272\,374\,985} \approx -0.0129367 \tag{38}$$

which is less than the value in Equation (31) because Banker has additional options while Player’s options are unchanged. The fully specified optimal strategy for Banker is given in Table 5.

This analysis extends to every positive integer d . We do not display the kernel, only the solution.

Table 5. Banker’s optimal move (final version) under Model B2 or B3 with d being a positive integer, indicated by D (draw), S (stand), or (S, D) (stand with probability $1 - q$, draw with probability q). In Model B2, q is as in Equations (39)–(42), and in Model B3, q is as in Equations (49)–(52).

Banker’s Two-Card Total	Player’s Third Card (\emptyset if Player Stands)											
	0	1	2	3	4	5	6	7	8	9	\emptyset	
0, 1, 2	D	D	D	D	D	D	D	D	D	D	D	D
3	D	D	D	D	D	D	D	D	S ¹	D	D	D
4	S	S ²	D	D	D	D	D	D	S	S	D	D
5	S	S	S	S	*	D	D	D	S	S	D	D
6	S	S	S	S	S	S	D ³	D	S	S	†	D
7	S	S	S	S	S	S	S	S	S	S	S	S

¹ Banker’s two-card total is 3 and Player’s third card is 8

(0, 3), (1, 2), (5, 8) S
 (4, 9) S if $d \geq 2$, D if $d = 1$
 (6, 7) S if $d \geq 1$ (Model B2)
 S if $d \geq 2$, D if $d = 1$ (Model B3)

² Banker’s two-card total is 4 and Player’s third card is 1

(0, 4), (1, 3), (2, 2), (5, 9) S
 (6, 8), (7, 7) S if $d \geq 3$, D if $d \leq 2$

* Banker’s two-card total is 5 and Player’s third card is 4

(0, 5), (7, 8) D if $d \geq 2$, S if $d = 1$
 (1, 4) S if $d \leq 7$, D if $d \geq 8$
 (2, 3) S if $d \leq 9$, D if $d \geq 10$ (Model B2)
 S if $d \leq 8$, D if $d \geq 9$ (Model B3)
 (6, 9) D if $d \geq 3$, S if $d \leq 2$ (Model B2)
 D if $d \geq 2$, S if $d = 1$ (Model B3)

³ Banker’s two-card total is 6 and Player’s third card is 6

(0, 6), (1, 5), (2, 4), (7, 9), (8, 8) D
 (3, 3) D if $d \geq 4$, S if $d \leq 3$

† Banker’s two-card total is 6 and Player stands

(0, 6) (S, D) (Model B2)
 (S, D) if $d \geq 2$, S if $d = 1$ (Model B3)
 (1, 5), (2, 4), (3, 3), (7, 9) S
 (8, 8) D (Model B2)
 D if $d \geq 2$, (S, D) if $d = 1$ (Model B3)

Theorem 6. Under Model B2 with d being a positive integer, the following Player and Banker strategies are uniquely optimal. Player draws on a two-card total of 5 with probability

$$p_d = \frac{(8d - 1)(12d - 1)(24d - 1)}{2d(1408d^2 - 220d + 9)}, \quad d \geq 1$$

and Banker draws on (0, 6), when Player stands, with probability

$$q_1 = \frac{290\,383}{450\,072}, \quad q_2 = \frac{2\,591\,845}{4\,119\,192}, \quad q_3 = \frac{9\,294\,089}{14\,521\,368} \tag{39}$$

$$q_d = \frac{368\,640d^4 - 68\,624d^3 - 2168d^2 + 981d - 48}{8d(52d - 5)(1408d^2 - 220d + 9)}, \quad 4 \leq d \leq 7 \tag{40}$$

$$q_d = \frac{367\,616d^4 - 67\,728d^3 - 2416d^2 + 1015d - 51}{8d(52d - 5)(1408d^2 - 220d + 9)}, \quad d = 8, 9 \tag{41}$$

or

$$q_d = \frac{366\,592d^4 - 67\,344d^3 - 2456d^2 + 1017d - 51}{8d(52d - 5)(1408d^2 - 220d + 9)}, \quad d \geq 10 \tag{42}$$

The value of the game (to Player) is

$$v_1 = -\frac{22\,932\,137}{1\,666\,583\,100}, \quad v_2 = -\frac{8\,220\,886\,553}{620\,866\,384\,425}, \quad v_3 = -\frac{210\,084\,639\,838}{16\,053\,072\,820\,785}$$

$$v_d = -32d(11\,125\,325\,824d^7 - 4\,182\,669\,312d^6 + 615\,333\,888d^5 - 43\,467\,904d^4 + 1\,329\,008d^3 + 5040d^2 - 1551d + 39)/[(1408d^2 - 220d + 9)(52d)_6], \quad 4 \leq d \leq 7$$

$$v_d = -32d(11\,129\,683\,968d^7 - 4\,218\,739\,712d^6 + 635\,681\,024d^5 - 47\,725\,760d^4 + 1\,738\,944d^3 - 14\,344d^2 - 1093d + 33)/[(1408d^2 - 220d + 9)(52d)_6], \quad d = 8, 9$$

or

$$v_d = -32d(11\,134\,042\,112d^7 - 4\,259\,389\,440d^6 + 648\,152\,320d^5 - 49\,007\,232d^4 + 1\,788\,256d^3 - 14\,816d^2 - 1089d + 33)/[(1408d^2 - 220d + 9)(52d)_6], \quad d \geq 10 \tag{43}$$

The fully specified optimal strategy for Banker is given in Table 5.

Remark. This is a slightly stronger statement than that of Downton and Lockwood [8].³ The reason for having different formulas for q_d and v_d in the six cases $d = 1, d = 2, d = 3, 4 \leq d \leq 7, d = 8, 9,$ and $d \geq 10$ is that Banker’s two pure strategies that determine the kernel do not vary if $4 \leq d \leq 7,$ if $d = 8, 9,$ or if $d \geq 10.$ This is a consequence of Table 5, which comes largely from Table 2a,b of [8].

³ They assumed $d \leq 8$ and rounded results to four decimal places, but they also allowed $d = \frac{1}{2},$ which we do not. Their Table 2b, which graphically represents the dependence on d in Banker’s optimal strategy, contains three ambiguities. Specifically, for $(j_1, j_2, k) = (6, 8, 1)$ and $d = 2,$ for $(j_1, j_2, k) = (2, 3, 4)$ and $d = 8,$ and for $(j_1, j_2, k) = (6, 9, 4)$ and $d = 2,$ it is uncertain whether D or S was intended. We have confirmed that D, S, and S, respectively, were intended in these cases. Furthermore, their table seems to suggest that D applies when $(j_1, j_2, k) = (2, 3, 4)$ and $d = 9,$ which is incorrect.

Proof. We apply Lemma 2 for $d = 1, 2, \dots, 19$. Only $(3, 3, 6)$ belongs to T_{10} (if $d \leq 10$). For each choice of d , the program runtime is about 15 seconds. For $d \geq 20$, the 22 points $p(l)$ ($l \in T_{01} \cup T_{10}$), relabeled as p_1, p_2, \dots, p_{22} , at which Player’s expected payoff is evaluated satisfy $p_{16} < p_{12} < p_{15} < p_{13} < p_{14} < p_5 < p_4 < p_1 < p_3 < p_2 < p_{22} < p_{17} < p_{21} < p_{18} < p_{19} < p_{20} < p_{11} < p_{10} < p_8 < p_7 < p_9 < p_6$, so the algorithm applies with a variable d . For $d \geq 20$, $V(p)$ is maximized at p_{17} (or $p(0, 6, \emptyset)$). At p_{22} (or $p(8, 8, \emptyset)$), for example,

$$V(p_{22}) = -32d(2783510528d^7 - 1188571136d^6 + 203128704d^5 - 16568896d^4 + 596408d^3 - 16158d^2 + 1855d - 93) / [(352d^2 - 71d + 4)(52d)_6]$$

and this is less than $V(p_{17})$ (see Equation (43)) for all $d \geq 20$, though the difference tends to 0 as $d \rightarrow \infty$. □

8. Solution under Model B3

We have reduced the game to a $2^5 \times 2^{n_d}$ matrix game, where $18 \leq n_d \leq 23$. The fact that Table 1 is identical under Models B2 and B3 suggests the existence of a 2×2 kernel under Model B3 whose column labels are the same as those of the 2×2 kernel under Model B2 as described in Section 7. Then the resulting $2^5 \times 2$ matrix game will of course have a 2×2 kernel, which is easy to find by graphical or other methods, and it will remain to confirm that this kernel corresponds to a solution of the $2^5 \times 2^{n_d}$ matrix game.

As we will see, this approach works for all positive integers d except 1, 2, and 9. These last three cases can be treated separately.

Let us begin with the case $d = 6$. Here the rows of the 32×2 payoff matrix are labeled from 0 to 31 as in Section 4, with the 5-bit binary form of the row number specifying the strategy (1 indicates D and 0 indicates S, in the five cases $(0, 5)$, $(1, 4)$, $(2, 3)$, $(6, 9)$, and $(7, 8)$). For example, row 19 (binary 10011) corresponds to drawing on $(0, 5)$, standing on $(1, 4)$ and $(2, 3)$, and drawing on $(6, 9)$ and $(7, 8)$.

We could label the two columns in a similar way but with binary strings of length 18 corresponding to the asterisks in Table 1, in the specific order $(0, 3, 9)$, $(1, 2, 9)$, $(4, 9, 9)$, $(5, 8, 9)$, $(6, 7, 9)$, $(2, 2, 1)$, $(6, 8, 1)$, $(7, 7, 1)$, $(0, 5, 4)$, $(6, 9, 4)$, $(7, 8, 4)$, $(3, 3, 6)$, $(0, 6, \emptyset)$, $(1, 5, \emptyset)$, $(2, 4, \emptyset)$, $(3, 3, \emptyset)$, $(7, 9, \emptyset)$, $(8, 8, \emptyset)$, reading the string left to right. In that case, the two columns would be labeled

$$111110001111000001 \quad \text{and} \quad 111110001111000001 \tag{44}$$

The kernel is easily found to be given by rows 19 and 27, so it is equal to

$$\begin{matrix} \text{B: S on } (0, 6), \emptyset & \text{B: D on } (0, 6), \emptyset \\ \text{P: S on } (1, 4) & \left(\begin{matrix} -19769569403 & -19425699931 \\ -19391857983 & -19783609631 \end{matrix} \right) / 1525814595305 \\ \text{P: D on } (1, 4) & \end{matrix}$$

implying that the following Player and Banker strategies are optimal. Player draws on $(1, 4)$ with probability

$$p_6 = \frac{35003}{74880} \approx 0.467455 \tag{45}$$

and Banker draws on (0, 6), when Player stands, with probability

$$q_6 = \frac{18\,885\,571}{36\,781\,056} \approx 0.513459 \tag{46}$$

The value of the game (to Player) is

$$v_6 = -\frac{73\,356\,216\,203\,119}{5\,712\,649\,844\,821\,920} \approx -0.0128410 \tag{47}$$

which is greater than the value in Equation (38) because Player has additional options while Banker’s options are unchanged. The fully specified optimal strategies for Player and Banker are given in Tables 5 and 6.

Table 6. Player’s optimal move under Model B3, indicated by D (draw), S (stand), or (S, D) (stand with probability 1 − p, draw with probability p). Here p is as in Equation (48).

Player’s Two-Card Hand	Optimal Move
(0, 5), (6, 9), (7, 8)	D
(1, 4)	(S, D)
(2, 3)	S

This analysis extends to every positive integer d. We do not display the kernel, only the solution.

Theorem 7. Under Model B3 with the number of decks being a positive integer d, the following Player and Banker strategies are optimal. Player draws on (1, 4) with probability

$$p_1 = \frac{1}{19}, \quad p_d = \frac{(12d - 1)(16d^2 - 14d + 1)}{32d^2(11d - 1)}, \quad d \geq 2 \tag{48}$$

and Banker draws on

$$\begin{cases} (8, 8) & \text{if } d = 1 \\ (0, 6) & \text{if } d \geq 2 \end{cases}$$

when Player stands, with probability

$$q_1 = \frac{4519}{10\,716}, \quad q_2 = \frac{17\,431}{64\,512}, \quad q_3 = \frac{4\,425\,647}{11\,132\,928} \tag{49}$$

$$q_d = \frac{92\,160d^4 - 120\,128d^3 + 26\,336d^2 - 2000d + 47}{256d^2(11d - 1)(52d - 5)}, \quad 4 \leq d \leq 7 \tag{50}$$

$$q_8 = \frac{316\,815\,305}{585\,842\,688} \tag{51}$$

or

$$q_d = \frac{91\,648d^4 - 119\,488d^3 + 26\,032d^2 - 1932d + 41}{256d^2(11d - 1)(52d - 5)}, \quad d \geq 9 \tag{52}$$

The value of the game (to Player) is

$$v_1 = -\frac{3\,439\,451}{25\,482\,800}, \quad v_2 = -\frac{49\,424\,010\,137}{3\,823\,801\,581\,600}, \quad v_3 = -\frac{31\,717\,439\,249}{2\,461\,444\,457\,472}$$

$$v_d = -2(1\,390\,665\,728\,d^7 - 491\,115\,520\,d^6 + 50\,698\,240\,d^5 + 2\,428\,032\,d^4 - 990\,512\,d^3 + 89\,192\,d^2 - 3462\,d + 47)/[(11\,d - 1)(52\,d)_6], \quad 4 \leq d \leq 7$$

$$v_8 = -\frac{2\,789\,416\,947\,665\,657}{217\,430\,324\,984\,396\,160}$$

or

$$v_d = -2(1\,391\,755\,264\,d^7 - 500\,535\,296\,d^6 + 54\,174\,464\,d^5 + 1\,931\,136\,d^4 - 948\,816\,d^3 + 85\,792\,d^2 - 3238\,d + 41)/[(11\,d - 1)(52\,d)_6], \quad d \geq 9 \tag{53}$$

The fully specified optimal strategies for Player and Banker are given in Tables 5 and 6.

Remark. We suspect that the solution is unique but do not have a proof. Notice what happens as $d \rightarrow \infty$: $p_d \rightarrow 6/11$, $q_d \rightarrow 179/286$, and $v_d \rightarrow -679\,568/53\,094\,899$, so the optimal mixed strategies approach those of Table 4.

Proof. Consider the case $d = 6$. Let \mathbf{A} denote the $2^5 \times 2^{18}$ payoff matrix, and let $\mathbf{p} = (p^0, p^1, \dots, p^{31})$ and $\mathbf{q} = (q^0, q^1, \dots, q^{262\,143})$ denote the stated optimal strategies. Specifically, with p_6 as in Equation (45) and q_6 as in Equation (46),

$$p^{27} = 1 - p^{19} = p_6 \quad \text{and} \quad p^i = 0 \text{ if } i \neq 19, 27$$

and

$$q^{254\,945} = 1 - q^{254\,913} = q_6 \quad \text{and} \quad q^j = 0 \text{ if } j \neq 254\,913, 254\,945$$

(254 913 and 254 945 are the decimal forms of the binary numbers in Equation (44).) Then, with v_6 as in Equation (47), it suffices to show that

$$\mathbf{pA} \geq (v_6, v_6, \dots, v_6) \tag{54}$$

componentwise (2^{18} inequalities), and

$$\mathbf{Aq}^T \leq (v_6, v_6, \dots, v_6)^T \tag{55}$$

componentwise (2^5 inequalities). Note that it is not necessary to evaluate the 2^{23} entries of \mathbf{A} . Equation (54) involves only rows 19 and 27 of \mathbf{A} (when rows are labeled 0 to 31), while Equation (55) involves only columns 254 913 and 254 945 of \mathbf{A} (when columns are labeled 0 to 262 143). We have confirmed Equations (54) and (55) using a *Mathematica* program. The program is, however, unnecessarily time-consuming. A more efficient way to proceed is to use Lemma 2 to verify Equation (54), with rows 19 and 27 (of 0–31) being Player’s two pure strategies. (In that case the number of exceptional cases is 6, namely $(0, 6, \emptyset)$, $(1, 5, \emptyset)$, $(2, 4, \emptyset)$, $(3, 3, \emptyset)$, $(7, 9, \emptyset)$, $(8, 8, \emptyset)$, not 18. T_{10} is empty.)

Similar programs give analogous results for every positive integer d . When d is 1, 2, or 9 and we use the two Banker pure strategies whose mixture is optimal under Model B2, we find that Equation (54) fails. By determining which components of the vector inequality fail, we can propose and confirm the correct optimal strategies under Model B3. Specifically, if $d = 2$ or $d = 9$, then Equation (54) fails at

just two components, namely the two that determine Banker's optimal mixed strategy under Model B3. If $d = 1$, then Equation (54) fails at eight components, which include the two optimal ones. Some trial and error may be required in this case.

We hasten to add that, just as in the case $d = 6$, there is a more efficient way. The two Player pure strategies specified by the kernel do not vary with d and are rows 19 and 27 (of 0–31). We again apply Lemma 2 to establish Equation (54), saving time and avoiding the issue that occurred in the cases $d = 1, 2, 9$. We find that, for $d = 1, 2, \dots, 9$, $n_d = 7, 9, 8, 9, 6, 6, 6, 7, 7$, respectively, and, if $d \geq 10$, $n_d = 6$. In each case, T_{10} is empty.

More importantly, this last method allows us to treat the cases $d \geq 10$ simultaneously, similarly to what we did in Section 7. For $d \geq 10$, the six points $p(l)$ ($l \in \{(0, 6, \emptyset), (1, 5, \emptyset), (2, 4, \emptyset), (3, 3, \emptyset), (7, 9, \emptyset), (8, 8, \emptyset)\}$), relabeled as p_1, p_2, \dots, p_6 , at which Player's expected payoff is evaluated, satisfy $p_6 < p_1 < p_5 < p_3 < p_2 < p_4$, so the algorithm applies with variable d . For $d \geq 10$, $V(p)$ is maximized at p_1 (or $p(0, 6, \emptyset)$). At p_6 (or $p(8, 8, \emptyset)$), for example,

$$V(p_6) = -8(695\,877\,632\,d^7 - 281\,198\,592\,d^6 + 34\,472\,064\,d^5 + 1\,177\,024\,d^4 - 901\,592\,d^3 + 119\,896\,d^2 - 6755\,d + 123)/[(22\,d - 3)(52\,d)_6]$$

and this is less than $V(p_1)$ (see Equation (53)) for all $d \geq 10$, though the difference tends to 0 as $d \rightarrow \infty$. \square

9. Summary

Baccara chemin de fer is a classical card game to which game theory is applicable. We consider six models, obtained by combining either Model A (cards are dealt with replacement) or Model B (cards are dealt without replacement from a d -deck shoe) with one of Model 1 (Player and Banker see two-card totals), Model 2 (Player sees two-card totals, Banker sees two-card composition), or Model 3 (Player and Banker see two-card composition). The first step in the analysis of these games is to reduce their size (2×2^{88} for Model 1, 2×2^{484} for Model 2, and $2^5 \times 2^{484}$ for Model 3) using strict dominance. Models A1, A2, and A3 are relatively simple and closely related. Model B1 is simple as well. Foster's algorithm can be applied to Model B2 in a straightforward way.

Model B3 is more challenging. Here we have a $2^5 \times 2^{n_d}$ matrix game, where $18 \leq n_d \leq 23$, after reduction. We guess that there is a 2×2 kernel whose column labels are the same as those found under Model B2. We then solve the resulting $2^5 \times 2$ game to obtain the two row labels of what we expect to be the kernel, and finally we confirm that this is indeed the kernel and we have a solution. It turns out that this method fails if $d = 1, 2$, or 9 . The two Banker pure strategies specified by the kernel are d -dependent, whereas the two Player pure strategies specified by the kernel are not. Thus, we can derive a solution for every positive integer d by applying Foster's algorithm to the 2×2^{n_d} matrix game whose row labels are those of the kernel.

As might be expected, the discrepancies between the values of the game under the various models are quite small. See Table 7.

Table 7. The value (to Player) of *baccara chemin de fer* as a function of the number of decks d under Models B1, B2, and B3, rounded to six significant digits. The values under Models A1, A2, and A3 are all equal to -0.0127991 .

d	Model B1	Model B2	Model B3
1	-0.0136323	-0.0137600	-0.0130413
2	-0.0131831	-0.0132410	-0.0129254
3	-0.0130391	-0.0130869	-0.0128857
4	-0.0129745	-0.0130113	-0.0128640
5	-0.0129390	-0.0129665	-0.0128505
6	-0.0129155	-0.0129367	-0.0128410
7	-0.0128988	-0.0129155	-0.0128341
8	-0.0128862	-0.0128998	-0.0128290
9	-0.0128765	-0.0128880	-0.0128254
10	-0.0128687	-0.0128791	-0.0128229
11	-0.0128624	-0.0128718	-0.0128209
12	-0.0128571	-0.0128657	-0.0128192

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Conflicts of Interest

The authors declare no conflict of interest.

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