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# CONTINUOUS TIME CONTESTS WITH PRIVATE INFORMATION

#### CHRISTIAN SEEL<sup>\*</sup> AND PHILIPP STRACK<sup>†</sup>

This paper introduces a class of contest models in which each player decides when to stop a privately observed Brownian motion with drift and incurs costs depending on his stopping time. The player who stops his process at the highest value wins a prize.

We prove existence and uniqueness of a Nash equilibrium outcome and derive the equilibrium distribution in closed form. As the variance tends to zero, the equilibrium outcome converges to the symmetric equilibrium of an all-pay auction. For two players and constant costs, each player's equilibrium profit decreases if the drift increases, the variance decreases, or the costs increase.

KEYWORDS: Contests, all-pay contests, silent timing games.

# 1. INTRODUCTION

There is a large economic literature on contests. Most of this literature analyzes static models (e.g., Tullock, 1980, Lazear and Rosen, 1981, Hillman and Samet, 1987, and Siegel, 2009, 2010). Most real-world contests, however, are conducted over time rather than being a one-shot interaction; for example, think about procurement contests, job promotion contests, political campaigns, or sports leagues. Thus, there are a few more recent papers which include the dynamic aspects of a contest in their model.

A dynamic contest model requires a choice of the informational assumptions. For instance, Harris and Vickers (1987) and Moscarini and Smith (2007) analyze stochastic and dynamic two-player contests with complete information. They characterize Markov perfect equilibria and find that the effort of the leader increases in his lead up to some threshold above which the laggard resigns. More general results on the existence of Markov perfect  $\epsilon$ -equilibria in two-player stopping games with complete information are provided in Laraki, Solan, and Vieille (2005) and Laraki and Solan (2013).

On the other hand, there is a literature which assumes that contestants cannot observe each other, so-called *silent timing games*; see, e.g., Karlin (1953) and Park and Smith (2008). In this literature, contestants can chose their stopping time, but they do not receive any new private information about their own process over time.

This paper introduces a new class of contest models without observability in which each player's strategies depend on their private information. If there is no uncertainty about the other players' progress, i.e., no private information, our model reduces to a silent

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timing game with one prize for the player who stops latest. Thus, this paper introduces private information into silent timing games.

Formally, we analyze an *n*-player contest in which each player decides when to stop a privately observed Brownian motion with drift. A strategy of a player is a bounded stopping time, i.e., the player has to stop almost surely before some deadline. As long as a player does not stop the process, he incurs flow costs of continuation. The player who stops his process at the highest value wins a prize.

The analysis proceeds as follows. Proposition 1 and Proposition 2 establish existence and uniqueness of the equilibrium distribution of the value of the stopped process. To do so, we first characterize the candidate equilibrium distribution uniquely up to its endpoints. In a second step, we use a result on Skorokhod embeddings to show that there exists a stopping strategy which leads to the candidate distribution.

In the paper, we do not fix the deadline exogenously, but the game ends once all players have stopped their processes. The equilibrium which we construct, however, remains an equilibrium if the contest ends at a sufficiently late, but fixed and predetermined deadline. Thus, we could use the more standard economic interpretation of a fixed deadline at which the results are evaluated if the deadline is sufficiently long. The bounded time restriction complicates the proof, since it drastically reduces the set of admissible stopping times.

By Proposition 3, an equilibrium stopping strategy has a simple structure, i.e., it is the first leaving time of a set. In economic terms, an agent stops whenever he either gets sufficiently optimistic or sufficiently pessimistic about his prospects in the contest, where the optimism and pessimism thresholds depend on time.

Having solved the model, we proceed by analyzing the shape of the equilibrium distribution. As uncertainty vanishes, the distribution converges to the symmetric equilibrium distribution of an all-pay auction by Proposition 4. On the one hand, the model offers a microfoundation for the use of all-pay auctions to scrutinize environments containing little uncertainty; on the other hand, it gives an equilibrium selection criterion for the equilibria of the symmetric all-pay auction analyzed in Baye, Kovenock, and de Vries (1996). Furthermore, this result serves as a benchmark to discuss how our predictions differ from all-pay models if volatility is strictly positive.

Most of the related literature focuses exclusively on two players and fixed cost functions. In the last part of the paper, we also restrict attention to two players and constant costs. This allows us to derive a tractable closed-form solution for the profits of each player. More precisely, by Proposition 6, profits increase as costs increase, volatility increases, or drift decreases. Hence, contestants prefer worse technologies for both players.

This economically novel comparative statics result has an intuitive explanation: an increase in the drift or a decrease in the variance of each player makes patience a more important factor for contest success compared to chance. Hence, in equilibrium, expected stopping times and expected costs increase. Since the equilibrium is symmetric, the winning probability of each player remains constant. Summing up, we obtain lower expected profits.

Our setting fits applications such as procurement contests or job promotion contests. In these applications, players choose their effort levels over time. The stopping decision measures the effort, where a later stopping time should be interpreted as a larger period of effort or a higher intensity of effort. Hence, in the model, a later stopping time also induces higher costs. Moreover, we assume that—on average—effort increases output. In the model, the average productivity is measured by the drift, i.e., the average rate at which the output increases.

In most procurement contests and some promotion promotion contests, the research is conducted in secrecy. Thus, contestants can only observe their own progress over time which is consistent with our model description. On the other hand, a model with complete information provides a better fit for applications such as sports leagues or political campaigns (if polls are available).

# Related Literature

Seel and Strack (2013) study the effect of relative performance pay on risk-taking behavior in financial markets. More precisely, fund managers decide when to stop investing in a risky asset. The manager who generates the highest profit gets a bonus. Each manager faces a liquidity constraint (maximal amount which they can loose) which drives the equilibrium construction. Seel and Strack (2013) construct an equilibrium in which managers invest in the risky asset although it has a negative expectation.

To focus on risk-taking behavior, Seel and Strack (2013) abstract from effort costs, which is the main decision variable in the classical contest literature. In contrast, all results in the present paper are driven by the trade-off between a higher expected cost and a higher probability of winning the contest.

Similar technical problems in characterizing the stopping region arise in Hopenhayn and Squintani (2011) who introduce private information into preemption games. They characterize the equilibrium as a solution to a free boundary problem, a technique which cannot be used in our setup. The solution method we develop is based on the Skorokhod embedding theorem. It provides a new approach to analyze timing games and, possibly, other stochastic games with private information. Moreover, as we show in an extension, the construction is not restricted to Brownian motion with drift, but can also be applied to other stochastic processes.

We proceed as follows. Section 2 sets up the model. In Section 3, we prove that an equilibrium exists and characterize the form of the equilibrium distribution uniquely.

Section 4 discusses the relation to all-pay auctions and derives the main comparative statics results. Section 5 concludes. All proofs are relegated to Appendix B.

## 2. THE MODEL

There are  $n \ge 2$  agents indexed by  $i \in \{1, 2, ..., n\} = N$  who face a stopping problem in continuous time. At each point in time  $t \in \mathbb{R}_+$ , agent *i* privately observes the realization of a stochastic process  $(X_t^i)_{t \in \mathbb{R}_+}$  with

$$X_t^i = \mu t + \sigma B_t^i \,.$$

The process  $X_i$  measures the (research) progress of a player who exerts effort. It consists of a deterministic and a stochastic component: the drift  $\mu \in \mathbb{R}_{++}$  is the common expected change of each process  $X_t^i$  per time, i.e.,  $\mathbb{E}(X_{t+\Delta}^i - X_t^i) = \mu\Delta$ . The random terms  $\sigma B_t^i$  are independent Brownian motions  $(B_t^i)_{i\in N}$  scaled by  $\sigma \in \mathbb{R}_{++}$ . Note that the starting value  $X_0^i$  is the same for each player and, without loss of generality, normalized to zero.

# 2.1. Strategies

A pure strategy of player *i* is a stopping time  $\tau^i$ . As each player only observes his own process, the decision whether to stop at time *t* only depends on the past realizations of his process  $(X_s^i)_{s \leq t}$ . Mathematically, the agents' stopping decision until time *t* has to be  $\mathcal{F}_t^i$ -measurable, where  $\mathcal{F}_t^i = \sigma(\{X_s^i : s < t\})$  is the sigma algebra induced by the possible observations of the process  $X_s^i$  before time *t*.

We require stopping times to be bounded by a real number  $T < \infty$  such that  $\tau^i < T$  almost surely.<sup>1</sup> Note that we do not fix the value of T in the description of the model, but only require T to be finite. However, provided T is long enough—we give a condition in Lemma 13—the equilibrium which we derive is also an equilibrium in a contest with a fixed deadline (see discussion after Proposition 3).

Although the unique equilibrium outcome of this model can be obtained in pure strategies, we incorporate mixing to make the results more general. To do so, we allow for random stopping decisions. More precisely, each agent *i* can choose an  $(\mathcal{F}_t^i)$ -adapted increasing process  $(\kappa_t^i)_{t\in\mathbb{R}_+}, \kappa_t^i \in [0,1]$  such that, for every  $t, \mathbb{P}(\tau^i \leq t | \mathcal{F}_t^i) = \kappa_t^i$ . A pure strategy is equivalent to a process  $\kappa_t^i$  that equals zero for all  $t < \tau$  and one otherwise.

## 2.2. Payoffs

The player who stops his process at the highest value wins a prize; ties are broken randomly. Until he stops, each player incurs flow costs  $c : \mathbb{R} \to \mathbb{R}_{++}$  which depend on the

<sup>&</sup>lt;sup>1</sup>The equilibrium we derive in the paper remains an equilibrium if we do not restrict agents to bounded stopping times, but only require stopping times to be finite, i.e.,  $\tau < \infty$ .



FIGURE 1.— An example for the game with two agents  $i \in \{1, 2\}$ . Time is depicted on the x-axis, while the value  $X_t^i$  (solid line) and total costs (dotted line) are depicted on the y-axis. The parameters are  $c(X_t^i) = \frac{1}{10} \int_0^t \exp(\frac{1}{10}X_s^i) ds$ ,  $\mu = 3$ ,  $\sigma = 1$ , T = 10.

current value of the process  $X_t^i$ , but not on the time t. We normalize the prize to one, since agents only care about the trade-off between winning probability and cost-prize ratio. The payoff  $\Pi^i$  is thus

$$\Pi^{i} = \frac{1}{k} \mathbf{1}_{\{X_{\tau^{i}}^{i} = \max_{j \in N} X_{\tau^{j}}^{j}\}} - \int_{0}^{\tau^{i}} c(X_{t}^{i}) \mathrm{d}t$$

where  $k = |\{i \in N : X_{\tau^i}^i = \max_{j \in N} X_{\tau^j}^j\}|$  is the number of agents who stop at the highest value. All agents maximize their expected profit  $\pi^i = \mathbb{E}(\Pi^i)$ . The cost function satisfies the following mild assumption:

ASSUMPTION 1 The cost function  $c : \mathbb{R} \to \mathbb{R}_{++}$  is continuous, positive, and bounded away from zero.

Figure 1 illustrates processes, stopping decisions, and corresponding costs in an example. Realizations of the processes until the stopping time are depicted in solid lines, while the corresponding total costs are depicted in dotted lines. In the example, player 1 stops before player 2 stops. Nonetheless, player 1 receives the prize, since she stops her process at a higher value than player 2.

#### 3. EQUILIBRIUM CONSTRUCTION

In this section, we first establish some necessary conditions on the distribution functions of the stopped processes in equilibrium. In a second step, we prove existence and uniqueness of the Nash equilibrium outcome and determine the equilibrium distributions depending on the cost function.

Every strategy of agent *i* induces a (potentially non-smooth) cumulative distribution function (cdf)  $F^i : \mathbb{R} \to [0,1]$  of his stopped process  $F^i(x) = \mathbb{P}(X^i_{\tau^i} \leq x)$ . Intuitively,  $F^i(x)$  is the probability that a player stops with an output less or equal to x.

Denote the endpoints of the support of the equilibrium distribution of player i by

$$\underline{x}^{i} = \inf\{x : F^{i}(x) > 0\} \overline{x}^{i} = \sup\{x : F^{i}(x) < 1\}.$$

Let  $\underline{x} = \max_{i \in N} \underline{x}^i$  and  $\overline{x} = \max_{i \in N} \overline{x}^i$ . In the next step, we establish a series of auxiliary results that are crucial to prove uniqueness of the equilibrium distribution.

LEMMA 1 No player places a mass point in the interior of the state space, i.e., for all i, for all  $x > \underline{x}$ :  $\mathbb{P}(X_{\tau^i}^i = x) = 0$ . At least one player has no mass at the left endpoint, i.e.,  $F^i(\underline{x}) = 0$ , for at least one player i.

We omit the proof of Lemma 1, since it is simply a specialization of the standard logic in static game theory with a continuous state space; see, e.g., Burdett and Judd (1983). Intuitively, in equilibrium, no player can place a mass point in the interior of the state space, since no other player would then stop slightly below the mass point.

Lemma 1 implies that the probability of a tie is zero. Thus, we can express the winning probability of player *i* if he stops at  $X_{\tau^i}^i = x$ , given the distributions of the other players, as

$$w^{i}(x) = \mathbb{P}(\max_{j \neq i} X^{j}_{\tau^{j}} \le x) = \prod_{j \neq i} F^{j}(x).$$

LEMMA 2 At least two players stop with positive probability on every interval  $I = (a, b) \subset [\underline{x}, \overline{x}].$ 

LEMMA 3 All players have the same equilibrium distribution function,  $F^i = F$ , for all  $i \in N$ .

Lemma 2 and 3 are the crucial steps for the uniqueness proof. In particular, Lemma 2 does not hold in static auction theory. Intuitively, the support of an equilibrium bid distribution in an auction is not necessary connected; see, for instance, Baye, Kovenock, and de Vries (1996), Theorem 1, who construct an equilibrium in which a player places a mass point at zero and randomizes his bid on some interval  $[x, \overline{x}]$ . In our case, every bounded time stopping strategy induces a distribution of the stopped process with a connected support. This property of bounded stopping strategies greatly facilitates the uniqueness proofs.

The following lemma is an immediate consequence from the symmetry of the equilibrium distribution (Lemma 3) and the fact that for at least one agent  $w^i(\underline{x}) = 0$  (Lemma 1).

LEMMA 4 Each player loses with certainty at  $\underline{x}$ , i.e.,  $w^i(\underline{x}) = 0$ , for all  $i \in N$ .

As players have symmetric distributions, we henceforth drop the superscript i. The previous lemmata imply that each player is indifferent between any stopping strategy on his support. The next lemma shows that the winning probability is twice differentiable on the support which allows us to use Itô's lemma.

LEMMA 5 The winning probability  $w : \mathbb{R} \to \mathbb{R}$  is twice differentiable for all  $x \in (\underline{x}, \overline{x})$ .

For any twice differentiable function  $\Phi : \mathbb{R} \to \mathbb{R}$  the evolution of the process  $(\Phi(X_t))_{t \in \mathbb{R}}$ is given by Itô's lemma

$$\mathrm{d}\Phi(X_t) = \left(\mu\Phi'(X_t) + \frac{\sigma^2}{2}\Phi''(X_t)\right)\mathrm{d}t + \sigma\Phi'(X_t)\mathrm{d}B_t\,.$$

As a consequence of Itô's lemma, the agent is indifferent if and only if for every point  $x \in (\underline{x}, \overline{x})$  the winning probability w satisfies the second order ordinary differential equation (ODE)

(1) 
$$c(x) = \mu w'(x) + \frac{\sigma^2}{2} w''(x)$$

The expected gain in winning probability from continuation equals the expected costs from continuation at every point of the support. As Eq. (1) is a second order ODE, we need two boundary conditions to determine w uniquely. One boundary condition is  $w(\underline{x}) = 0$  from Lemma 4. We determine the other one in the following lemma:

LEMMA 6 In equilibrium,  $w'(\underline{x}) = 0$ .

Intuitively, Lemma 6 makes sure that players have no incentive to continue at the left endpoint of their support.

Imposing the two boundary conditions, the solution to Eq. (1) is unique. To calculate it, we define  $\phi(x) = \exp(\frac{-2\mu x}{\sigma^2})$  as a solution to the homogeneous equation  $0 = \mu w'(x) + \frac{\sigma^2}{2} w''(x)$ . To solve the inhomogeneous equation, we apply the variation-ofconstants formula. We then use the two boundary conditions to calculate the unique candidate solution. Finally, we rearrange with Fubini's Theorem to get

$$w(x) = \begin{cases} 0 & \text{for } x < \underline{x} \\ \frac{1}{\mu} \int_{\underline{x}}^{x} c(z)(1 - \phi(x - z)) dz & \text{for } x \in [\underline{x}, \overline{x}] \\ 1 & \text{for } \overline{x} < x \,. \end{cases}$$

By symmetry of the equilibrium strategy, the function  $F : \mathbb{R} \to [0, 1]$  satisfies  $F(x) = \sqrt[n-1]{w(x)}$ . Consequently, the unique candidate for an equilibrium distribution is

(2) 
$$F(x) = \begin{cases} 0 & \text{for all } x \leq \underline{x} \\ \min\left\{1, \sqrt[n-1]{\frac{1}{\mu} \int_{\underline{x}}^{x} c(z)(1 - \phi(x - z)) \mathrm{d}z} \right\} & \text{for all } x > \underline{x} \end{cases}$$

In the next step, we verify that F is a cumulative distribution function.

LEMMA 7 For every left endpoint  $\underline{x}$ , the function F defined in Eq. (2) is nondecreasing and  $\lim_{x\to\infty} F(x) = 1$ .

The next result shows that there is unique candidate for the endpoints of each player's distribution:

PROPOSITION 1 There exists a unique pair 
$$(\underline{x}, \overline{x}) \in \mathbb{R}^2$$
 such that the distribution  
(3)  $F(x) = \begin{cases} 0 & \text{for all } x \leq \underline{x} \\ \sqrt[n-1]{\frac{1}{\mu} \int_{\underline{x}}^{\underline{x}} c(z)(1 - \phi(x - z)) \mathrm{d}z} & \text{for all } x \in (\underline{x}, \overline{x}) \\ 1 & \text{for all } x \geq \overline{x} \end{cases}$ 

is the unique candidate for an equilibrium distribution.

The next lemma shows that this condition is also sufficient.

LEMMA 8 Every strategy that induces the distribution F defined in Eq. (3) is an equilibrium strategy.

By construction of F, all agents are indifferent between all stopping strategies which stop inside the support  $[\underline{x}, \overline{x}]$ . As every agent wins with probability one at the right endpoint  $\overline{x}$ , it is strictly optimal to stop there. The condition  $F'(\underline{x}) = 0$  from Lemma 6 ensures that it is also optimal to stop at the left endpoint.

So far, we have verified that a bounded stopping time  $\tau \leq T < \infty$  is an equilibrium strategy if and only if it induces the distribution F, i.e.,  $F(z) = \mathbb{P}(X_{\tau} \leq z)$ . To show that the game has a Nash equilibrium, the existence of a bounded stopping time inducing Fremains to be established. The problem of finding a stopping time  $\tau$  such that a Brownian motion stopped at  $\tau$  has a given centered probability distribution F, i.e.,  $F \sim B_{\tau}$ , is known in the probability literature as the Skorokhod embedding problem (SEP). Since its initial formulation in Skorokhod (1961, 1965), many solutions have been derived; for a survey article, see Oblój (2004). The proof of the following proposition verifies a sufficient condition for a bounded time embedding from a recent mathematical paper.

**PROPOSITION 2** The game has a Nash equilibrium. In equilibrium, the strategy of each player induces the distribution F defined in Eq. (3).

The uniqueness result differs from related models such as symmetric all-pay auctions or symmetric silent timing games, in which there are usually multiple equilibria with different equilibrium distributions. In some of these equilibria only a subset of players is active, i.e., submits a positive bid in the auction or does not stop directly in the silent timing game. Here, the variance leads all players to be active in the game. Hence, we obtain a unique equilibrium distribution.

A conceptual improvement on the literature is the bounded time requirement  $\tau < T$ , which makes the equilibrium derivation applicable to contests with a fixed deadline. To see that this result is not trivial to obtain, note that for any fixed time horizon T, there exists a positive probability that  $X_t$  does not leave any interval [a, b] with  $a < X_0 < b$ . Hence, a simple construction through a mixture over cutoff strategies of the form

$$\tau_{a,b} = \inf\{t \in \mathbb{R}_+ : X_t \notin [a,b]\},\$$

cannot be used to implement F.

In the following, we show that a different simple strategy induces the equilibrium distribution. For this purpose, we define:

DEFINITION 1 (Barrier Stopping Time) A stopping time  $\tau$  is a barrier stopping time if it is of the form

$$\tau = \inf\{t \in \mathbb{R}_+ : (t, X_t) \in B\},\$$

where  $B \subset \mathbb{R}_+ \times \mathbb{R}$  is a set such that  $(t, x) \in B$  implies  $(t', x) \in B$  if t' > t.

Barrier stopping times are relatively simple, as they only condition on the value of the process and the time, but not on the history.

**PROPOSITION 3** There exists a bounded equilibrium stopping time which is a barrier stopping time.

An equilibrium in which each player uses a barrier stopping has a nice economic intuition: whenever a player gets sufficiently optimistic or sufficiently pessimistic about her prospects in the contest, she stops to exert effort. In Figure 2, we draw a sample path of the process and a barrier stopping time. The agent is sufficiently optimistic if she hits the upper threshold, while she is sufficiently pessimistic if she hits the lower threshold. Figure 2 also illustrates our result for contests with a fixed deadline: a barrier stopping time which induces the distribution defined in Eq. (3) stops almost surely before some finite point in time. If the deadline of the contest is greater or equal to this point in time, each player can thus induce the distribution defined in Eq. (3). Since the strategy space



FIGURE 2.— An example of a barrier stopping time.

is restricted by the deadline, the equilibrium derived for the larger strategy space without a fixed deadline remains an equilibrium in the game with a fixed, but sufficiently long deadline. For given parameters, it is possible to calculate an upper bound for the critical contest length using Lemma 13 in the appendix.

#### 4. EQUILIBRIUM ANALYSIS

#### 4.1. Convergence to the All-pay Auction

This subsection considers the relationship between the literature on all-pay auctions and our model for vanishing noise. We first establish an auxiliary result about the left endpoint:

LEMMA 9 If the noise vanishes, the left endpoint of the equilibrium distribution defined in Eq. (3) converges to zero, i.e.,  $\lim_{\sigma \to 0} \underline{x} = 0$ .

Taking the limit  $\sigma \to 0$ , the equilibrium distribution converges to

$$\lim_{\sigma \to 0} F(x) = \sqrt[n-1]{\frac{1}{\mu} \int_0^x c(z) \mathrm{d}z}.$$

In a static n-player all-pay auction, the symmetric equilibrium distribution is

$$F(x) = \sqrt[n-1]{\frac{x}{v}},$$

where x is the total outlay of a participant and v is her valuation; see, e.g., Hillman and Samet (1987). In our case, the total outlay depends on the flow costs at each point, the

speed of research  $\mu$ , and the stopping time  $\tau$ . More precisely, it is  $\int_0^x \frac{c(z)}{\mu} dz$ . The valuation v in the analysis of Hillman and Samet (1987) coincides with the prize p—which we have normalized to one—in our contest. This yields the following proposition:

**PROPOSITION 4** For vanishing noise, the equilibrium distribution converges to the symmetric equilibrium distribution of an all-pay auction with homogeneous valuations.

Thus, our model supports the use of all-pay auctions to analyze contests in which variance is negligible. Moreover, the symmetric all-pay auction has multiple equilibria for a full characterization see Baye, Kovenock, and de Vries (1996), Theorem 1. This paper offers an equilibrium selection criterion in favor of the symmetric equilibrium. Intuitively, all other equilibria of the symmetric all-pay auction include mass points at zero for some players, which is not possible in our model for any positive  $\sigma$  by Lemma 4.

#### 4.2. Comparative Statics and Rent Dispersion

Proposition 4 has linked all-pay auctions with complete information to our model for the case of vanishing noise. In the following, we scrutinize how predictions differ for positive noise. In a symmetric all-pay auction with complete information, agents make zero profits in equilibrium. This does not hold true in our model for any positive level of variance  $\sigma$ :

## **PROPOSITION 5** In equilibrium, all agents make strictly positive expected profits.

Intuitively, agents generate informational rents through the private information about their research progress. A similar result is known in the literature on all-pay auctions with incomplete information, see, e.g., Hillman and Riley (1989), Amann and Leininger (1996), Krishna and Morgan (1997), and Moldovanu and Sela (2001). In these models, each participant takes a draw from a common distribution. The outcome of each draw is private information and determines the effort cost or valuation. In contrast to this, private information about one's progress arrives continuously over time in our model.

# 4.3. Two Players and Constant Costs

We henceforth restrict attention to the case n = 2 and c(x) = c to get more explicit results. We require some additional notation. In particular, we denote the principal branch of the Lambert W-function by  $W_0: [-\frac{1}{e}, \infty) \to \mathbb{R}_+$ . This branch is implicitly defined on  $[-\frac{1}{e}, \infty)$  as the unique solution of  $x = W(x) \exp(W(x)), W \ge -1$ . Moreover, define  $h: \mathbb{R}_+ \to [0, 1]$  by

$$h(y) = \exp(-y - 1 - W_0(-\exp(-1 - y))).$$



FIGURE 3.— This figure shows the equilibrium profit  $\pi$  of the agents on the *y*-Axes for n = 2 constant cost-functions  $c(x) = c \in \mathbb{R}_{++}$  with  $y = \frac{2\mu^2}{c\sigma^2}$  on the *x*-Axes.

The next lemma pins down the left and right endpoints of the support.

LEMMA 10 The left and right endpoint are  

$$\underline{x} = \frac{\sigma^2}{2\mu} \left( 2\log(1 - h(\frac{2\mu^2}{c\sigma^2})) - \log(\frac{4\mu^2}{c\sigma^2}) \right)$$

$$\overline{x} = \frac{\sigma^2}{2\mu} \left( 2\log(1 - h(\frac{2\mu^2}{c\sigma^2})) - \log(\frac{4\mu^2}{c\sigma^2}) - \log(h(\frac{2\mu^2}{c\sigma^2})) \right).$$

The next lemma derives a closed-form solution of the profits  $\pi$  of each player.

LEMMA 11 The equilibrium profit of each player depends only on the ratio  $y = \frac{2\mu^2}{c\sigma^2}$ . It is given by

$$\pi = \frac{(1 - h(y))^2}{2y^2} - \frac{2\log(1 - h(y)) - \log(2y) - 1}{y}$$

For an illustration of the above formula, see Figure 3. The previous lemma enables us to establish the main comparative static result of this paper.

**PROPOSITION 6** The equilibrium profit of each player increases if costs increase, variance increases, or drift decreases.

To get an intuition for the result, we decompose the term  $\frac{2\mu^2}{c\sigma^2}$ , which determines the equilibrium profit of the players, into two parts:

$$\frac{2\mu^2}{c\sigma^2} = \underbrace{\frac{\mu}{c}}_{\text{Productivity}} \times \underbrace{\frac{2\mu}{\sigma^2}}_{\text{Signal to noise ratio}}.$$

The first term  $\frac{\mu}{c}$  is a deterministic measure of productivity per unit of time. On the other

hand, the second term  $\frac{2\mu}{\sigma^2}$  measures the impact of randomness  $\sigma^2$  on the final outcome. As the drift of both competitors increases or the variance decreases, patience becomes more critical to winning than chance. Hence, the expected stopping time in equilibrium increases and expected total costs increase. As each player's winning probability remains  $\frac{1}{2}$  by symmetry, profits decrease. In summary, participants prefer to have worse—more costly, more random, or less productive—technologies.

#### 4.4. Extension to Other Stochastic Processes

Up to now, we have derived our results for Brownian motion with drift. However, it is straightforward to extend them to any process  $Z_t$  such that  $Z_t$  is the result of a strictly monotone transformation  $\rho : \mathbb{R} \to \mathbb{R}$  of Brownian motion with drift  $X_t$ , i.e.,  $Z_t = \rho(X_t)$ . First, observe that the success process  $Z_t$  is only payoff relevant as an ordinal variable, as agent *i* receives a prize if

$$Z^{i}_{\tau^{i}} > \max_{j \neq i} Z^{j}_{\tau^{j}} \Leftrightarrow \rho^{-1}(Z^{i}_{\tau^{i}}) > \max_{j \neq i} \rho^{-1}(Z^{j}_{\tau^{j}}) \Leftrightarrow X^{i}_{\tau^{i}} > \max_{j \neq i} X^{j}_{\tau^{j}}$$

As costs c depend on the value of the process  $Z_t$ , we define the cost function  $\tilde{c}(x) = c(\rho(x))$ for the process  $X_t$ . Let  $\tilde{F}$  be the equilibrium distribution for the game with a Brownian motion with drift  $X_t$  and cost  $\tilde{c}$ , i.e.,  $X_{\tau^i}^i \sim \tilde{F}$ . The equilibrium distribution of the original game is now given by  $F(x) = \tilde{F}(\rho^{-1}(x))$ , i.e.,  $Z_{\tau^i}^i \sim F$ . For example, this generalization includes geometric Brownian motion with drift  $(Z_t = \exp(X_t))$ .

#### 5. CONCLUSION AND DISCUSSION

In this paper, we have introduced a new class of contest models in continuous time with private information. Under mild assumptions on the cost function, a Nash equilibrium exists. We have characterized the unique equilibrium distribution of the stopped process. Moreover, there exists a bounded stopping strategy—the first leaving time of a set—which leads to this distribution. For two players and constant costs, each player prefers mutually higher research costs, worse technologies, and higher uncertainty. This economically novel feature results from chance becoming more important for contest success than patience. From a technical perspective, we have introduced a method to construct equilibria in continuous time games. This methodological contribution should enhance future research on timing games and other stochastic games with private information.

The derivation of the equilibrium relies on the Markovian structure of each player's optimization problem. It remains an open question how to deal with non-Markovian settings.

#### 6. APPENDIX A: PRELIMINARY RESULTS

Before we prove the main results from the text in Appendix B, we introduce two technical auxiliary results and a remark which turn out to be useful for the main proofs.

Recall the definition  $\phi(x) = \exp(\frac{-2\mu x}{\sigma^2})$ . The next lemma derives a necessary condition for a distribution F to be the outcome of a strategy  $\tau$ .

LEMMA 12 If  $\tau \leq T < \infty$  is a bounded stopping time that induces the absolutely continuous distribution F, i.e.,  $F(z) = \mathbb{P}(X_{\tau} \leq z)$ , then

(4) 
$$1 = \int_{\underline{x}}^{\overline{x}} \phi(x) F'(x) \mathrm{d}x.$$

**PROOF:** Observe that  $(\phi(X_t))_{t \in \mathbb{R}_+}$  is a martingale. Hence, by Doob's optional stopping theorem, for any bounded stopping time  $\tau$ ,

$$1 = \phi(X_0) = \mathbb{E}[\phi(X_\tau)] = \int_{\underline{x}}^{\overline{x}} \phi(x) F'(x) \mathrm{d}x \,. \qquad Q.E.D.$$

Moreover, we state a result from Ankirchner and Strack (2011). In addition to the assumption stated in the next lemma, Ankirchner and Strack (2011) assume that the condition in Lemma 12 holds. They define  $g(x) = F^{-1}(\Phi(x))$ , where  $\Phi(x) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{x} \exp(\frac{z^2}{2}) dz$ is the density function of the normal distribution.

LEMMA 13 (Ankirchner and Strack, 2011, Theorem 2) Suppose that g is Lipschitzcontinuous with Lipschitz constant  $\sqrt{T}$ . Then the distribution F can be embedded in  $X_t = \mu t + B_t$ , with a stopping time that stops almost surely before T.

Let us denote the first leaving time of the interval [a, b] of agent i after time t by

$$\tau_{(a,b)}^{t,i} = \inf\{s \ge t : X_s^i \notin (a,b)\}.$$

**REMARK 1** The continuation strategy  $\tau_{(a,b)}^{t,i}$  is strictly better than stopping at x for agent i iff

$$\mathbb{E}\left(w^{i}(X^{i}_{\tau^{t,i}_{(a,b)}}) - \int_{t}^{\tau^{t,i}_{(a,b)}} c(X^{i}_{s}) \mathrm{d}s \,|\, X^{i}_{t} = x\right) > w^{i}(x)$$

If  $\tau_{(a,b)}^{t,i}$  is a profitable continuation strategy, the continuation strategy  $\tau_{a,b}^{t,i} \wedge T$  is also

profitable for a sufficiently long, but finite time horizon T, since

$$\lim_{T \to \infty} \mathbb{E} \left( w^i (X^i_{\tau^{t,i}_{(a,b)} \wedge T}) - \int_t^{\tau^{(i,i)}_{(a,b)} \wedge T} c(X^i_s) \mathrm{d}s \,|\, X^i_t = x \right)$$
$$= \mathbb{E} \left( w^i (X^i_{\tau^{t,i}_{(a,b)}}) - \int_t^{\tau^{t,i}_{(a,b)}} c(X^i_s) \mathrm{d}s \,|\, X^i_t = x \right).$$

Thus, in the next proofs, it suffices to show that  $\tau_{(a,b)}^{t,i}$  is a profitable deviation to rule out equilibrium candidates.

#### 7. APPENDIX B

We are now ready to prove our main results:

PROOF OF LEMMA 2: As players have to use bounded time stopping strategies, each player *i* stops with positive probability on every subinterval of  $[\underline{x}^i, \overline{x}^i]$  (for a formal proof of this statement, see Ankirchner and Strack, 2011). Hence, it suffices to show that at least two players have  $\overline{x}$  as their right endpoint. Assume, by contradiction, only player *i* has  $\overline{x}$  as his right endpoint. Denote  $\overline{x}^{-i} = \max_{j \neq i} \overline{x}^j$ . Then, by continuity of  $F^j$ , player *i* strictly prefers to stop at  $\overline{x}^{-i}$ , which yields him the maximal possible winning probability of 1 without any additional costs. This contradicts the optimality of a strategy, which stops at  $\overline{x}^i > \overline{x}^{-i}$ . Q.E.D.

PROOF OF LEMMA 3:: Step 1:  $F^{i}(x) \leq F^{j}(x)$  if and only if  $w^{j}(x) \leq w^{i}(x)$  for all  $x > \underline{x}$ .

By definition,  $w^i(x)F^i(x) = w^j(x)F^j(x)$ . Hence,  $\frac{w^i(x)}{w^j(x)} = \frac{F^j(x)}{F^i(x)}$ .

**Step 2:**  $w^i(\underline{x}^j) \ge w^j(\underline{x}^j)$  implies  $F^i(x) \le F^j(x)$  for almost all x.

**Step 2.1:** We first prove  $w^j(x) \leq w^i(x)$  for all  $x \in (\underline{x}^j, \overline{x}^j) \cap (\underline{x}^i, \overline{x}^i)$ . For all  $x \in (\underline{x}^j, \overline{x}^j) \cap (\underline{x}^i, \overline{x}^i)$  player *i* weakly prefers stopping at *x* to the continuation strategy  $\tau^{t,i}_{(\underline{x}^j, \overline{x}^j)}$ 

$$\mathbb{E}\left(w^{i}(X^{i}_{\tau^{t,i}_{(\underline{x}^{j},\overline{x}^{j})}}) - \int_{0}^{\tau^{t,i}_{(\underline{x}^{j},\overline{x}^{j})}} c(X^{i}_{s}) \mathrm{d}s \,|\, X^{i}_{t} = x\right) \leq w^{i}(x) \,.$$

As  $\underline{x}^{j}$  and  $\overline{x}^{j}$  are in the support of player j the continuation strategy  $\tau_{(\underline{x}^{j},\overline{x}^{j})}^{t,j}$  is optimal for every point  $x \in (\underline{x}^{j}, \overline{x}^{j})$ , which yields

$$\mathbb{E}\left(w^{j}(X^{j}_{\tau^{t,j}_{(\underline{x}^{j},\overline{x}^{j})}}) - \int_{0}^{\tau^{t,j}_{(\underline{x}^{j},\overline{x}^{j})}} c(X^{j}_{s}) \mathrm{d}s \,|\, X^{j}_{t} = x\right) \geq w^{j}(x) \,.$$

By assumption,  $w^i(\underline{x}^j) \ge w^j(\underline{x}^j)$ . Moreover,  $F^i(\overline{x}) \le 1 = F^j(\overline{x}^j)$  which yields by Step 1  $w^i(\overline{x}^j) \ge w^j(\overline{x}^j)$ . Hence  $w^j(X^j_{\tau^{t,j}_{(\underline{x},\overline{x})}}) \le w^i(X^i_{\tau^{t,i}_{(\underline{x}^j,\overline{x}^j)}})$  and as

$$\mathbb{E}\left(\int_{0}^{\tau_{(\underline{x}^{j},\overline{x}^{j})}^{t,i}} c(X_{s}^{i}) \mathrm{d}s \,|\, X_{t}^{i} = x\right) = \mathbb{E}\left(\int_{0}^{\tau_{(\underline{x}^{j},\overline{x}^{j})}^{t,j}} c(X_{s}^{j}) \mathrm{d}s \,|\, X_{t}^{j} = x\right)$$

the above equations imply  $w^j(x) \leq w^i(x)$  for all  $x \in (\underline{x}^j, \overline{x}^j) \cap (\underline{x}^i, \overline{x}^i)$ .

**Step 2.2:** In this step, we prove  $w^i(x) \ge w^j(x)$  for all  $x \ge \overline{x}^i \wedge \overline{x}^j$ . By Step 2.1,  $F^i(x) \le F^j(x)$  for all  $x \in (\underline{x}^j, \overline{x}^j) \cap (\underline{x}^i, \overline{x}^i)$ . As  $1 = F^i(\overline{x}^i) \le F^j(\overline{x}^i)$  it follows that  $\overline{x}^j \le \overline{x}^i$ . As  $F^i(x) \le 1 = F^j(x)$  it follows by Step 1 that  $w^j(x) \le w^i(x)$  for all  $x \ge \overline{x}^j$ .

**Step 2.3:** In the final step, we prove  $w^i(\underline{x}^j) \ge w^j(\underline{x}^j)$  for all  $x \le \underline{x}^i \lor \underline{x}^j$ . First let  $\underline{x}^i < \underline{x}^j$ . As no agent places a mass point in the interior it follows that  $F^j(\underline{x}^j) = 0$  and thus  $w^i(\underline{x}^j) = 0$ . Using the result of the first step it follows that for all  $x \le \underline{x}^j$  $w^j(\underline{x}^j) \le \lim_{x \searrow \underline{x}^j} w^j(x) \le \lim_{x \searrow \underline{x}^j} w^i(x) = 0$ . Hence,  $w^i(x) = w^j(x) = 0$  for all  $x \le \underline{x}^j$ . Let  $\underline{x}^j \le \underline{x}^i$ . As  $F^i(x) = 0 \le F^j(x)$  for all  $x < \underline{x}^i$  it follows from Step 1 that  $w^j(x) \le w^i(x)$  for all  $x < \underline{x}^i$ .

Hence,  $w^{j}(x) \leq w^{i}(x)$  almost everywhere and by Step 1  $F^{i} \leq F^{j}$  almost everywhere.

Step 3:  $F^i = F^j$  for all i, j.

If  $\underline{x}^i = \underline{x}^j$  it follows by Step 2 that either  $F^i \leq F^j$  or  $F^i \geq F^j$  almost everywhere. Otherwise assume, without loss of generality,  $\underline{x}^j < \underline{x}^i$ . Thus,  $0 = w^j(\underline{x}^j) \leq w^i(\underline{x}^j)$  and hence, by Step 2,  $F^i \leq F^j$  almost everywhere. Hence, w.l.o.g., we assume for the rest of the proof  $F^i \leq F^j$  almost everywhere. If  $F^i < F^j$  on a set with positive Lebesgue measure it follows that

$$\int_{\mathbb{R}_+} (1 - F^j(x))\phi'(x) \mathrm{d}x < \int_{\mathbb{R}_+} (1 - F^i(x))\phi'(x) \mathrm{d}x.$$

This violates the feasibility condition (Lemma 12). Thus  $F^i = F^j$  almost everywhere. By continuity of  $F^i$  and  $F^j$ , it follows that  $F^i = F^j$ . Q.E.D.

PROOF OF LEMMA 5: Let  $q : \mathbb{R} \to \mathbb{R}$  be the solution to the differential equation  $\mu q'(x) - \frac{\sigma^2}{2}q''(x) = c(x)$  with  $q(\underline{x}) = 0$  and  $q'(\underline{x}) = 0$ . Then,

$$\mathbb{E}\left(\int_{t}^{\min\{\tau_{(\underline{x},\overline{x})},K\}} c(X_s) \mathrm{d}s \,|\, X_t = x\right) = \mathbb{E}(q(X_{\min\{\tau_{(\underline{x},\overline{x})},K\}}) - q(X_t) \,|\, X_t = x)$$

As the agent is indifferent inside the support the strategy  $\min\{\tau_{(\underline{x},\overline{x})}, K\}$  is an optimal continuation strategy. Thus, the winning probability w equals

$$w(x) = \lim_{K \to \infty} \mathbb{E} \left( w(X_{\min\{\tau_{(\underline{x},\overline{x})},K\}}) - \int_{t}^{\min\{\tau_{(\underline{x},\overline{x})},K\}} c(X_{s}) \mathrm{d}s \,|\, X_{t} = x \right)$$
$$= \lim_{K \to \infty} \mathbb{E} \left( w(X_{\min\{\tau_{(\underline{x},\overline{x})},K\}}) - q(X_{\min\{\tau_{(\underline{x},\overline{x})},K\}}) + q(X_{t}) \,|\, X_{t} = x \right)$$
$$= \mathbb{P}(X_{\tau_{(\underline{x},\overline{x})}} = \overline{x} \,|\, X_{t} = x)(1 - q(\overline{x})) + q(x) \,.$$

The result follows as q(x) and  $\mathbb{P}(X_{\tau_{(\underline{x},\overline{x})}} = \overline{x} \mid X_t = x)$  are twice differentiable in x. Q.E.D.

PROOF OF LEMMA 6: By definition, w(x) = 0, for all  $x \leq \underline{x}$ . Hence, the left derivative  $\partial_-w(\underline{x})$  is zero. It remains to prove that the right derivative  $\partial_+w(\underline{x})$  is also zero. For a given  $w: \mathbb{R} \to \mathbb{R}_+$ , let  $\Psi: \mathbb{R} \to \mathbb{R}$  be the unique function that satisfies the second order ordinary differential equation  $c(x) = \mu \Psi'(x) + \frac{\sigma^2}{2} \Psi''(x)$  with boundary conditions  $\Psi(\underline{x}) = 0$  and  $\Psi'(\underline{x}) = \partial_+w(\underline{x})$ . As  $\Psi'(\underline{x}) > 0$ , there exists a point  $\hat{x} < \underline{x}$  such that  $\Psi(\hat{x}) < 0 = w(\hat{x})$ . Consider the strategy  $S = \min\{\tau_{\hat{x},\overline{x}}, 1\}$  that stops when either the point  $\hat{x}$  or  $\overline{x}$  is reached or at time 1. As  $w(\hat{x}) > \Psi(\hat{x})$ , it follows that  $\mathbb{E}(w(X_S)) > \mathbb{E}(\Psi(X_S))$ . Thus,

$$\mathbb{E}(w(X_S) - \int_0^S c(X_t^i) \mathrm{d}t) > \mathbb{E}(\Psi(X_S) - \int_0^S c(X_t^i) \mathrm{d}t) \,.$$

Note that, by Itô's lemma, the process  $\Psi(X_t^i) - \int_0^t c(X_s^i) ds$  is a martingale. By Doob's optional sampling theorem, agent *i* is indifferent between the equilibrium strategy  $\tau$  and the bounded time strategy *S*, i.e.,

$$\mathbb{E}(\Psi(X_S) - \int_0^S c(X_t^i) \mathrm{d}t) = \mathbb{E}(\Psi(X_\tau) - \int_0^\tau c(X_t^i) \mathrm{d}t) = \mathbb{E}(w(X_\tau) - \int_0^\tau c(X_t^i) \mathrm{d}t).$$

The last step follows because w(x) and  $\Psi(x)$  coincide for all  $x \in (\underline{x}, \overline{x})$ . Consequently, the strategy S is a profitable deviation, which contradicts the equilibrium assumption. Q.E.D.

PROOF OF LEMMA 7: By construction of F,  $F(\underline{x}) = 0$ . Clearly, F is increasing on  $(\underline{x}, \overline{x})$ , as the derivative with respect to x,

$$F'(x) = \frac{F(x)^{2-n}}{(n-1)} \left(\frac{2}{\sigma^2} \int_{\underline{x}}^x c(z)\phi(x-z)\mathrm{d}z\right) \,,$$

is greater than zero for all  $x > \underline{x}$ . It remains to show that there exists an  $x > \underline{x}$  such that

F(x) = 1. We can bound F from below using a lower bound on the derivative of  $F^{n-1}$ :

$$F(x)^{n-1} = \frac{1}{\mu} \int_{\underline{x}}^{x} c(z)(1 - \phi(x - z)) dz$$
  
$$\geq \frac{c}{\mu} \left( x - \underline{x} - \frac{\sigma^{2}}{2\mu} (1 - \phi(x - \underline{x})) \right) \geq \frac{c}{\mu} \left( x - \underline{x} - \frac{\sigma^{2}}{2\mu} \right)$$

By Assumption 1,  $\underline{c} > 0$ . Continuity of F implies that there exists a unique point  $\overline{x} > \underline{x}$  such that  $F(\overline{x}) = 1$ . Q.E.D.

PROOF OF PROPOSITION 1: Uniqueness: As F is absolutely continuous, the right endpoint  $\overline{x}$  satisfies  $(a) : 1 = \int_{\underline{x}}^{\overline{x}} F'(x; \underline{x}, \overline{x}) dx$ . By Lemma 7, there exists a unique  $\overline{x}$ for every  $\underline{x}$  such that (a) is satisfied. Lemma 12 states that any feasible distribution satisfies  $(b) : 1 = \int_{\underline{x}}^{\overline{x}} F'(x; \underline{x}) \phi(x) dx$ . We show that intersection of the set of solutions to equation (a) and to equation (b) consists of a single point. Since  $F'(x; \underline{x}, \overline{x})$  is independent of  $\overline{x}$ , we henceforth drop the dependency in our notation. By the implicit function theorem, the set of solutions to (a) satisfies

(5) 
$$\frac{\partial \overline{x}}{\partial \underline{x}} = -\frac{-\overbrace{F'(\underline{x};\underline{x})}^{=0} + \int_{\underline{x}}^{\overline{x}} \frac{\partial}{\partial \underline{x}} F'(x;\underline{x}) dx}{F'(\overline{x};\underline{x})} = -\frac{\int_{\underline{x}}^{\overline{x}} \frac{\partial}{\partial \underline{x}} F'(x;\underline{x}) dx}{F'(\overline{x};\underline{x})}$$

Applying the implicit function theorem to (b) gives us

$$(6) \qquad \frac{\partial \overline{x}}{\partial \underline{x}} = -\frac{\overbrace{F'(\underline{x};\underline{x})}^{=0} + \int_{\underline{x}}^{\overline{x}} \frac{\partial}{\partial \underline{x}} F'(x;\underline{x})\phi(x)dx}{F'(\overline{x};\underline{x})\phi(\overline{x})} = -\frac{\int_{\underline{x}}^{\overline{x}} \frac{\partial}{\partial \underline{x}} F'(x;\underline{x}) \phi(\overline{x}-\overline{x})dx}{F'(\overline{x};\underline{x})} dx$$
$$< -\frac{\int_{\underline{x}}^{\overline{x}} \frac{\partial}{\partial \underline{x}} F'(x;\underline{x})dx}{F'(\overline{x};\underline{x})}.$$

The last inequality follows from  $\frac{\partial}{\partial \underline{x}}F'(x;\underline{x}) \ge 0$ . As (5) > (6), the solution sets to equation (a) and (b) intersect at most once. Thus, in equilibrium, both the left and the right endpoint are unique.

**Existence:** We have shown in Lemma 7 that, for every  $\underline{x}$ , there exists a unique  $\overline{x}$  such that  $F(\cdot; \underline{x}, \overline{x})$  is a distribution function. Furthermore, rewriting the last inequality in the proof of Lemma 7, we get  $\overline{x} \leq \underline{x} + \frac{\mu}{c} + \frac{\sigma^2}{2\mu}$  and thus  $\underline{x} \to -\infty \Rightarrow \overline{x} \to -\infty$ . Consider a left endpoint  $\underline{x} \geq 0$  and a right endpoint  $\overline{x}$  such that (a) is satisfied. Then,  $\phi(x) < 1$  for all  $x \in [\underline{x}, \overline{x}]$  and hence  $\mathbb{E}(\phi(x)) < 1$  for  $x \sim F(\cdot, \underline{x}, \overline{x})$ . Now consider a left endpoint  $\underline{x} < 0$  and a right endpoint  $\overline{x} < 0$  such that (a) is satisfied. Then,  $\phi(x) > 1$  for all  $x \in [\underline{x}, \overline{x}]$  and hence  $\mathbb{E}(\phi(x)) > 1$  for  $x \sim F(\cdot, \underline{x}, \overline{x})$ . Now consider a left endpoint  $\underline{x} < 0$  and a right endpoint  $\overline{x} < 0$  such that (a) is satisfied. Then,  $\phi(x) > 1$  for all  $x \in [\underline{x}, \overline{x}]$  and hence  $\mathbb{E}(\phi(x)) > 1$  for  $x \sim F(\cdot, \underline{x}, \overline{x})$ . Equation (b) is equivalent to  $\mathbb{E}(\phi(x)) = 1$ . By continuity and the Intermediate Value Theorem, there exists an  $\underline{x}, \overline{x}$  such (a) and (b) are satisfied.

PROOF OF LEMMA 8: Define  $\Psi$  as the unique solution to Eq. (1) with the boundary conditions  $\Psi(\underline{x}) = 0$  and  $\Psi'(\underline{x}) = 0$ . By construction, the process  $\Psi(X_t^i) - \int_0^t c(X_s^i) ds$  is a martingale and  $\Psi(x) = w(x)$  for all  $x \in [\underline{x}, \overline{x}]$ . As  $\Psi'(x) < 0$  for  $x < \underline{x}$  and  $\Psi'(x) > 0$ for  $x > \overline{x}$ ,  $\Psi(x) > w(x)$  for all  $x \notin [\underline{x}, \overline{x}]$ . For every stopping time S, we use Itô's Lemma to calculate the expected value

$$\mathbb{E}\left(w(X_S) - \int_0^S c(X_t) dt\right) \leq \mathbb{E}\left(\Psi(X_S) - \int_0^S c(X_t) dt\right)$$
  
=  $\Psi(X_0) = w(X_0) = \mathbb{E}(w(X_\tau)).$ 

The last equality results, as each agent is indifferent to stop immediately with the expected payoff  $w(X_0)$  or to play the equilibrium strategy with the expected payoff  $\mathbb{E}(w(X_{\tau}))$ . The inequality implies that any possible deviation strategy obtains a weakly lower payoff than the equilibrium candidate. Q.E.D.

PROOF OF PROPOSITION 2: The function  $\Phi$  is Lipschitz continuous with constant  $\frac{1}{\sqrt{2\pi}}$ . Consequently, it suffices to prove Lipschitz continuity of  $F^{-1}$  to get the Lipschitz continuity of  $F^{-1} \circ \Phi$ . The density f is

$$f(x) = \frac{F(x)^{-n+2}}{n-1} \frac{2}{\sigma^2} \int_{\underline{x}}^x c(z)\phi(x-z)dz = \frac{F(x)^{-n+2}}{n-1} \frac{2}{\sigma^2} \left( \int_{\underline{x}}^x c(z)dz - \mu F(x)^{n-1} \right).$$

As f(x) > 0 for all  $x > \underline{x}$ , it suffices to show Lipschitz continuity of  $F^{-1}$  at 0. We substitute  $x = F^{-1}(y)$  to get

$$(f \circ F^{-1})(y) \ge \frac{1}{n-1} \frac{2}{\sigma^2} \left( y^{2-n} (\min_{z \in [\underline{x}, \overline{x}]} c(z)) (F^{-1}(y) - F^{-1}(0)) - \mu y \right).$$

We define  $\underline{c} = \inf_{z \in [\underline{x}, \overline{x}]} c(z)$ . Rearranging with respect to  $F^{-1}(y) - F^{-1}(0)$  gives

$$F^{-1}(y) - F^{-1}(0) \le \left(\frac{(n-1)\sigma^2}{2}(f \circ F^{-1})(y) + \mu y\right) \frac{y^{n-2}}{\underline{c}} \le \left(\frac{(n-1)\sigma^2}{2}f(\overline{x}) + \mu\right) \frac{y^{n-2}}{\underline{c}} + \frac{y^{n-2}}{\underline$$

This proves the Lipschitz continuity of  $F^{-1}$  for n > 2. Note that for two agents n = 2 the function  $F^{-1}$  is not Lipschitz continuous as  $f(\underline{x}) = 0$ . However, we show in the following paragraph that  $F^{-1} \circ \Phi$  is Lipschitz continuous for n = 2.

$$F(x) = \int_{\underline{x}}^{x} \frac{c(z)}{\mu} (1 - \phi(x - z) dz \le \left( \sup_{z \in [\underline{x}, \overline{x}]} \frac{c(z)}{\mu} \right) \left( x - \underline{x} - \frac{\sigma^2}{2\mu} (1 - \phi(x - \underline{x})) \right)$$

We define  $\overline{c} = \sup_{z \in [\underline{x}, \overline{x}]} c(z)$ . A second order Taylor expansion around  $\underline{x}$  yields that, for an open ball around  $\underline{x}$  and  $\underline{x} < x$ , we have the following upper bound

$$x - \underline{x} - \frac{\sigma^2}{2\mu} (1 - \phi(x - \underline{x})) \le \frac{2\mu}{\sigma^2} (1 - \phi(x - \underline{x}))^2$$

For an open ball around  $\underline{x}$ , we get an upper bound on  $F(x) \leq \frac{2\overline{c}}{\sigma^2}(1-\phi(x-\underline{x}))^2$  and hence the following estimate

$$1 - \phi(x - \underline{x}) \ge \sqrt{\frac{\sigma^2}{2\overline{c}}F(x)}$$

We use this estimate to obtain a lower bound on f depending only on F

$$f(x) = \frac{2}{\sigma^2} \left( \int_{\underline{x}}^x c(z)\phi(x-z) dz \right) \ge \frac{2\underline{c}}{\sigma^2} \left( \frac{\sigma^2}{2\mu} (1 - \phi(x-\underline{x})) \right) \ge \frac{\underline{c}}{\mu} \sqrt{\frac{\sigma^2}{2\overline{c}}} F(x) \,.$$

Consequently, there exists an  $\epsilon > \underline{x}$  such that, for all  $x \in [\underline{x}, \epsilon)$ , we have an upper bound on  $\frac{(\phi \circ \Phi^{-1} \circ F)(x)}{f(x)}$ . Taking the limit  $x \to \underline{x}$  yields

$$\lim_{x \to \underline{x}} \frac{(\phi \circ \Phi^{-1} \circ F)(x)}{f(x)} \leq \lim_{x \to \underline{x}} \frac{(\phi \circ \Phi^{-1} \circ F)(x)}{\frac{\underline{c}}{\mu} \sqrt{\frac{\sigma^2}{2\overline{c}} F(x)}} \leq \sqrt{\frac{2\overline{c}\mu^2}{\underline{c}^2 \sigma^2}} \lim_{y \to 0} \frac{(\phi \circ \Phi^{-1})(y)}{\sqrt{y}} = 0.$$

Hence, we have shown that  $F^{-1} \circ \Phi$  is Lipschitz continuous. Thus, there exists a bounded time stopping strategy (for an explicit construction, see Ankirchner and Strack, 2011, Theorem 2) which embeds F. Q.E.D.

PROOF OF PROPOSITION 3: The stopping strategy constructed in Ankirchner and Strack (2011) is not a barrier stopping time. However, Rost (1976) shows in his Theorem 1 that if there exists a stopping time that embeds the distribution F, there also exists a barrier stopping time which embeds F. Furthermore, by Theorem 2 in Rost (1976), this first leaving time of a set minimizes

$$\mathbb{E}(\tau - \min\{\tau, t\}) = \mathbb{P}(\tau \ge t)\mathbb{E}(\tau - t \mid \tau \ge t)$$

over the set of all stopping times that embed F for all  $t \in \mathbb{R}_+$ . As the stopping time constructed in Ankirchner and Strack (2011) is bounded by T, the first leaving time of a set constructed in Rost (1976) is also bounded by T. Q.E.D.

PROOF OF LEMMA 9: For any bounded stopping time, for any  $\sigma > 0$ , feasibility implies that  $\underline{x} \leq 0$ . By contradiction, assume there exists a constant  $\epsilon$  such that  $\underline{x} \leq \epsilon < 0$  for some sequence  $(\sigma_k)_{k \in \mathbb{N}}$  with  $\lim_{k \to \infty} \sigma_k = 0$ . Then F' is bounded away from zero by

$$F'(x) = \frac{F(x)^{2-n}}{n-1} \frac{2}{\sigma^2} \int_{\underline{x}}^x c(z)\phi(x-z)dz \ge \frac{1}{n-1} \frac{2}{\sigma^2} \int_{\epsilon}^x c(z)\phi(x-z)dz$$
$$\ge \frac{\underline{c}}{\mu(n-1)} (1-\phi(x-\epsilon)).$$

For every point x < 0,  $\lim_{\sigma_k \to 0} \phi(x) = \infty$ . Thus,

$$\lim_{\sigma_k \to 0} \int_{\underline{x}}^0 F'(x)\phi(x) \mathrm{d}x > 1 \,,$$

which contradicts feasibility, because

$$\int_{\underline{x}}^{0} F'(x)\phi(x)\mathrm{d}x \le \int_{\underline{x}}^{\overline{x}} F'(x)\phi(x)\mathrm{d}x = 1. \qquad Q.E.D.$$

PROOF OF PROPOSITION 5: In equilibrium, agents are indifferent between stopping immediately and the equilibrium strategy. Their expected profit is thus w(0), which is strictly positive as  $\underline{x} < 0$ . Q.E.D.

PROOF OF LEMMA 10 AND LEMMA 11: Define  $\Delta = \overline{x} - \underline{x}$ . Rearranging the density condition  $1 = F(\overline{x}) = \frac{c}{\mu} [\Delta - \frac{\sigma^2}{2\mu} (1 - \phi(\Delta))]$  yields

$$\exp\left(\frac{-2\mu\Delta}{\sigma^2}\right) = -\frac{2\mu}{\sigma^2}\left[\Delta - \left(\frac{\mu}{c} + \frac{\sigma^2}{2\mu}\right)\right]$$

The solution to the transcendental algebraic equation  $e^{-a\Delta} = b(\Delta - d)$  is  $\Delta = d + \frac{1}{a}W_0(\frac{ae^{-ad}}{b})$ , where  $W_0 : [-\frac{1}{e}, \infty) \to \mathbb{R}_+$  is the principal branch of the Lambert W-function. This branch is implicitly defined on  $[-\frac{1}{e}, \infty)$  as the unique solution of  $x = W(x) \exp(W(x)), W \ge -1$ . Hence,

$$\Delta = \frac{\mu}{c} + \frac{\sigma^2}{2\mu} \left[ 1 + W_0 \left( -\exp(-1 - \frac{2\mu^2}{c\sigma^2}) \right) \right]$$

and

$$\phi(\Delta) = \exp\left(\frac{-2\mu^2}{c\sigma^2} - 1 - W_0(-\exp(-1 - \frac{2\mu^2}{c\sigma^2}))\right)$$
  
=  $\exp\left(-1 - y - W_0(-\exp(-1 - y))\right) = h(y)$ .

Note that  $\phi(\Delta)$  only depends on  $y = \frac{2\mu^2}{c\sigma^2}$ . As  $\Delta \ge 0$  it follows that  $h(y) = \phi(\Delta) \in [0, 1]$ . Moreover,  $h : \mathbb{R}_+ \to [0, 1]$  is strictly decreasing in y, as  $W_0$  and exp are strictly increasing functions. For constant costs, the feasibility condition from Lemma 12 reduces to

$$1 = \int_{\underline{x}}^{\overline{x}} F'(x)\phi(x)dx = \frac{c\sigma^2}{2\mu^2} \left[\frac{1}{2}\phi(\underline{x}) + \frac{1}{2}\phi(2\overline{x} - \underline{x}) - \phi(\overline{x})\right].$$

Dividing by  $\phi(\underline{x})$  gives

$$\begin{aligned} \phi(-\underline{x}) &= \frac{c\sigma^2}{2\mu^2} \left[ \frac{1}{2} + \frac{1}{2}\phi(\Delta)^2 - \phi(\Delta) \right] &= \frac{1}{y} \left[ \frac{1}{2} + \frac{1}{2}h(y)^2 - h(y) \right] \\ &= \frac{1}{2y} \left( 1 - h(y) \right)^2 = g(y) \end{aligned}$$

As  $\underline{x} \leq 0$  it follows that  $g(y) = \phi(-\underline{x}) \in [0, 1]$ . The derivative of g is given by

$$g'(y) = -\frac{1}{2y^2}(1 - h(y))^2 + \frac{2}{2y}(1 - h(y))h'(y)$$

As h is a strictly decreasing function from  $\mathbb{R}_+$  to [0,1] it follows that  $g:\mathbb{R}_+\to [0,1]$  is strictly decreasing in y. We calculate  $\underline{x}$  as

$$\underline{x} = -\phi^{-1}(\phi(-\underline{x})) = \frac{\sigma^2}{2\mu} \log(g(y)) = \frac{\sigma^2}{2\mu} \log\left(\frac{1}{2y} \left(1 - h(y)\right)^2\right)$$
$$= -\frac{\sigma^2}{2\mu} [2\log\left(1 - h(y)\right) - \log(2y)] \,.$$

Simple algebraic transformations yield the expression for  $\overline{x}$  in Lemma 10. We plug in  $\underline{x} = \frac{\sigma^2}{2\mu} \log(g(y))$  and  $\phi(-\underline{x}) = g(y)$  in the distribution for the constant cost case to get:

$$\begin{split} F(0) &= w(0) = \frac{c}{\mu} \left[ -\underline{x} - \frac{\sigma^2}{2\mu} (1 - \phi(-\underline{x})) \right] = \frac{c}{\mu} \left[ -\frac{\sigma^2}{2\mu} \log(g(y)) - \frac{\sigma^2}{2\mu} (1 - g(y)) \right] \\ &= \frac{1}{y} [g(y) - \log(g(y)) - 1] \,. \end{split}$$

Hence, the value of F(0) depends on the value of the fraction  $y = \frac{2\mu^2}{c\sigma^2}$  in the above way, which completes the proof of Lemma 11. Q.E.D.

PROOF OF PROPOSITION 6: By Lemma 11, it suffices to show that the profit  $\pi = w(X_0) = w(0) = F(0)$  is increasing in y. Consider the following expression from the previous proof:

$$F(0) = \frac{1}{y} [g(y) - \log(g(y)) - 1] .$$

Note  $\frac{1}{y}$  is positive and  $g(y) - \log(g(y)) - 1$  is positive as  $F(0) \ge 0$ . The function  $x - \log(x)$  is increasing in x. Hence,  $g(y) - \log(g(y)) - 1$  is decreasing in y, because g(y) is decreasing in y. As  $\frac{1}{y}$  is also decreasing in y, the product  $\frac{1}{y}[g(y) - \log(g(y)) - 1]$  of two decreasing positive functions is decreasing in y. Q.E.D.

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